

WEBASSEMBLY illustrated

exploring some mental models and implementations

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WIP

Rev. 0.01.0

NOTE

- Please refer to the official documents in detail.
- This information is based on "WebAssembly Specification Release 1.0 (Draft, last updated Oct 31, 2018)".
- This information is current as of Nov, 2018.
Still work in progress.

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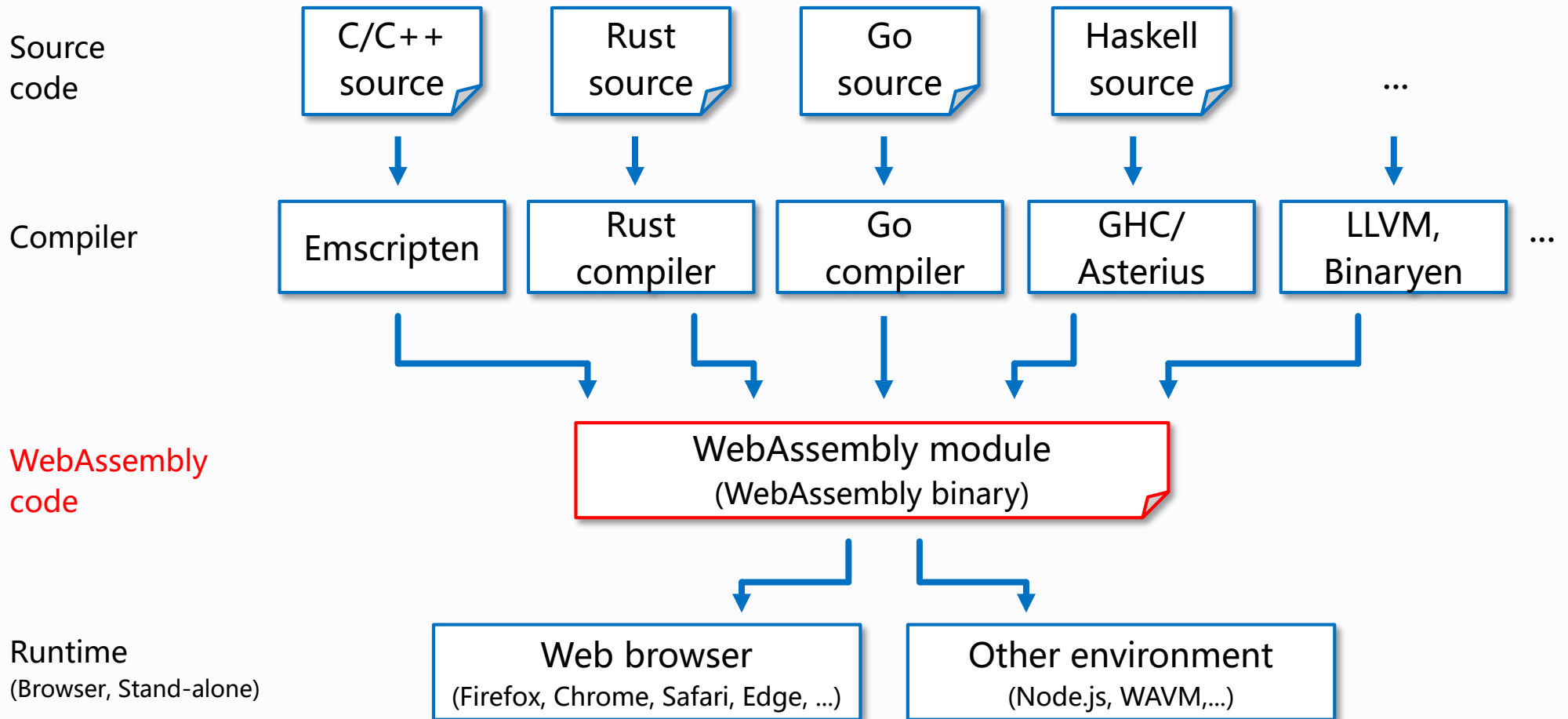
References

1. Introduction

1. Introduction

Overview

WebAssembly is a code format



WebAssembly is a safe, portable, low-level code format.

WebAssembly code

Text format

Binary format

syntactic sugar

```
(module
  (func (export "add7")
    (param $x i64)
    (result i64)
    i64.add
    (get_local $x)
    i64.const 7))))
```

core syntax

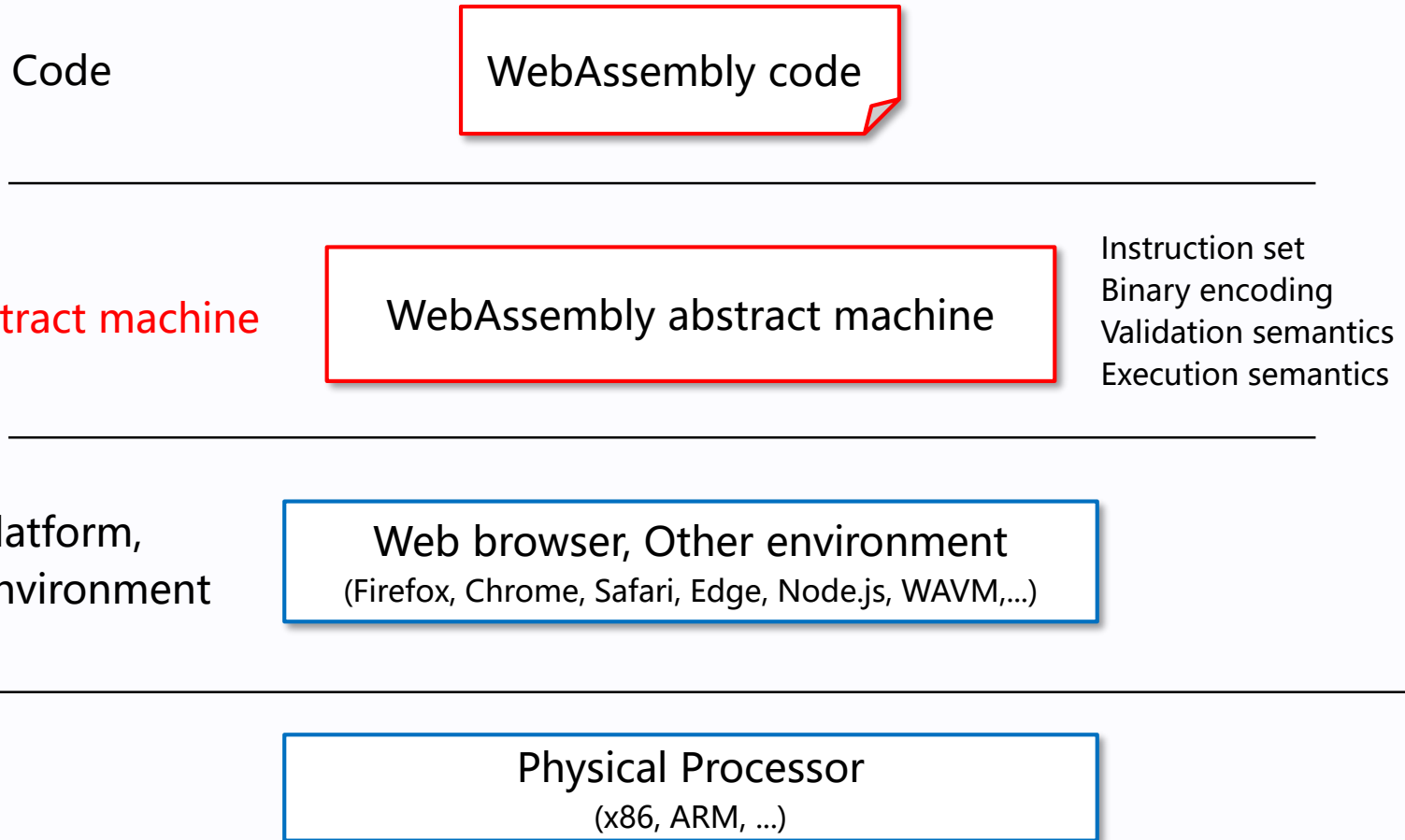
```
(module
  (type
    (func (param i64) (result i64)))
  (func (type 0)
    (param i64) (result i64)
    get_local 0
    i64.const 7
    i64.add)
  (export "add7" (func 0)))
```

```
0x0061736d010000 ...
```

WebAssembly encodes a low-level, assembly-like programming language.

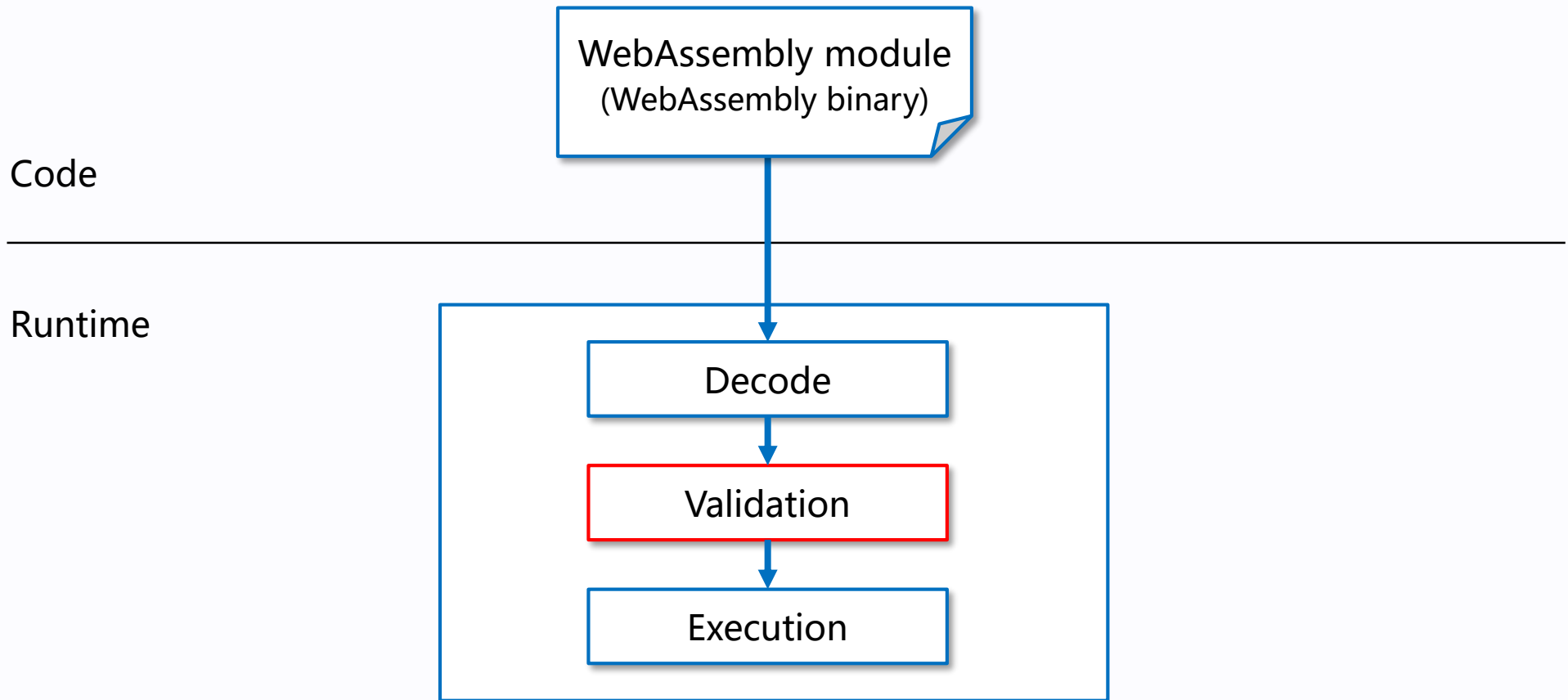
WebAssembly has multiple concrete representations.
(its text format and the binary format.)

Abstract machine is defined



WebAssembly is a virtual instruction set architecture (virtual ISA). Execution behavior is defined in terms of an abstract machine.

Validation



Validation checks that a WebAssembly module is well-formed.

Validity is defined by a type system.

The type system of WebAssembly is sound, implying both type safety and memory safety with respect to the WebAssembly semantics.

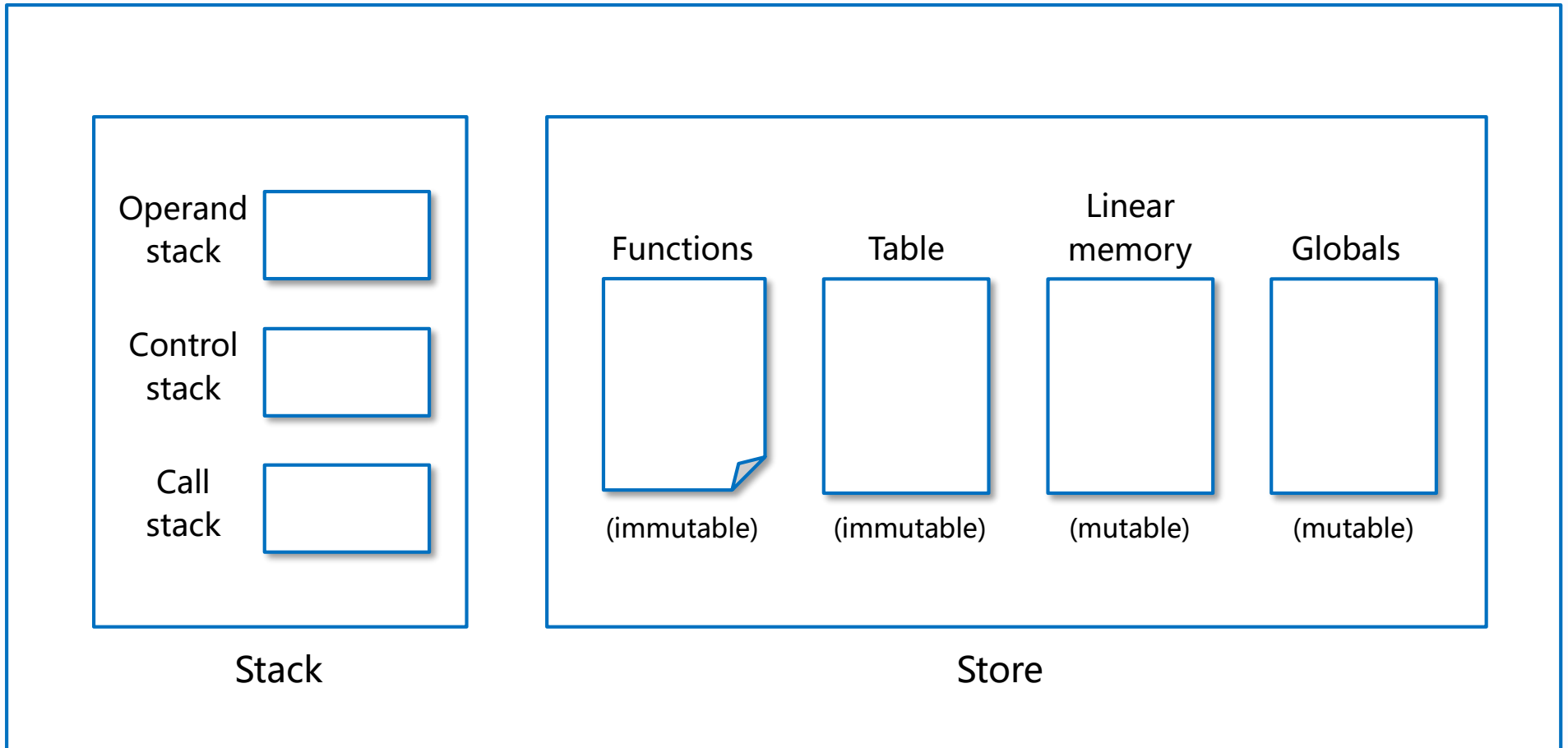
2. WebAssembly abstract machine

2. WebAssembly abstract machine

Abstract machine

WebAssembly abstract machine

WebAssembly abstract machine

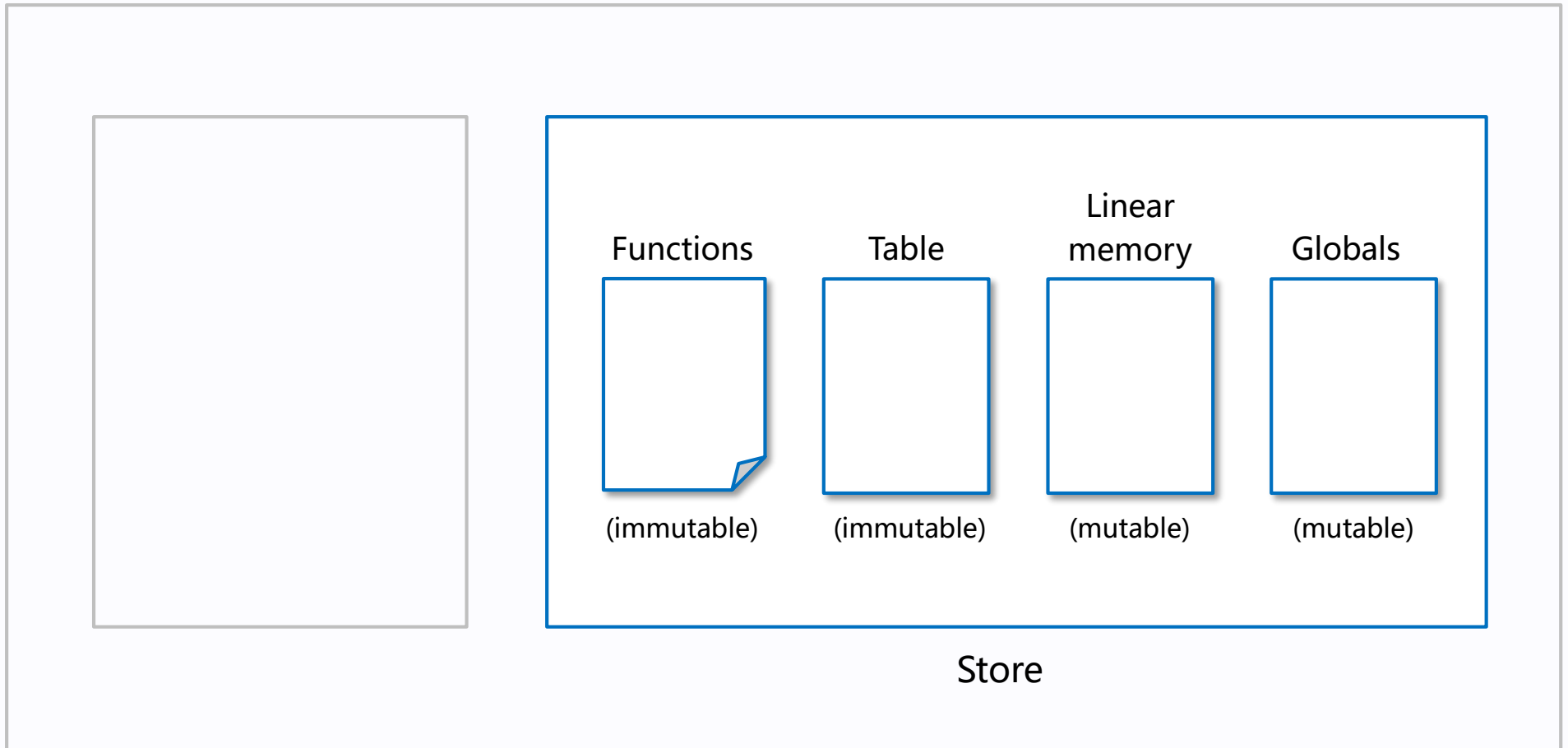


WebAssembly abstract machine is based on a stack machine.
The abstract machine includes a store and an implicit stack.

2. WebAssembly abstract machine

Store

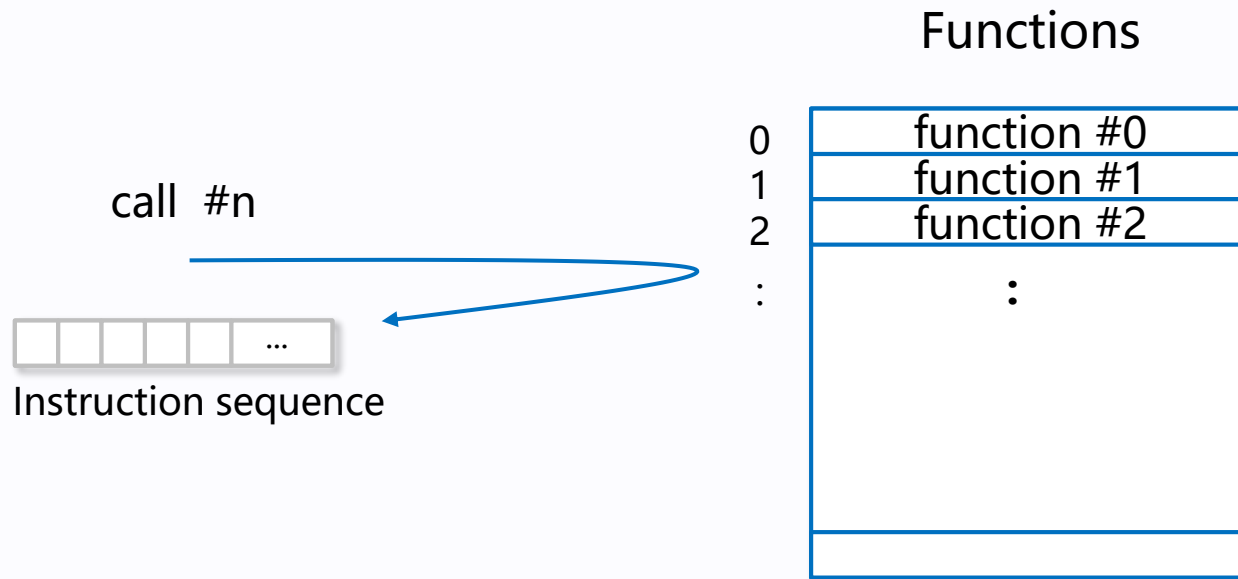
Store



The store represents all global state.

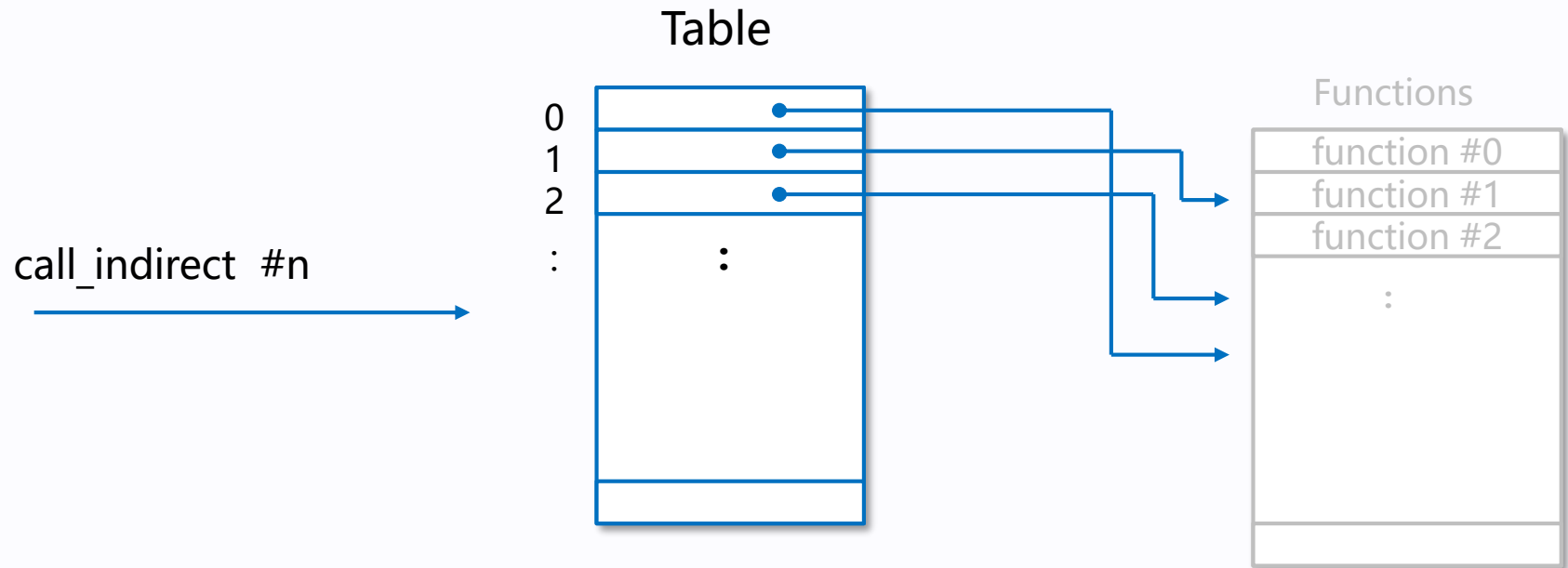
The store have been allocated during the life time of the abstract machine.

Functions



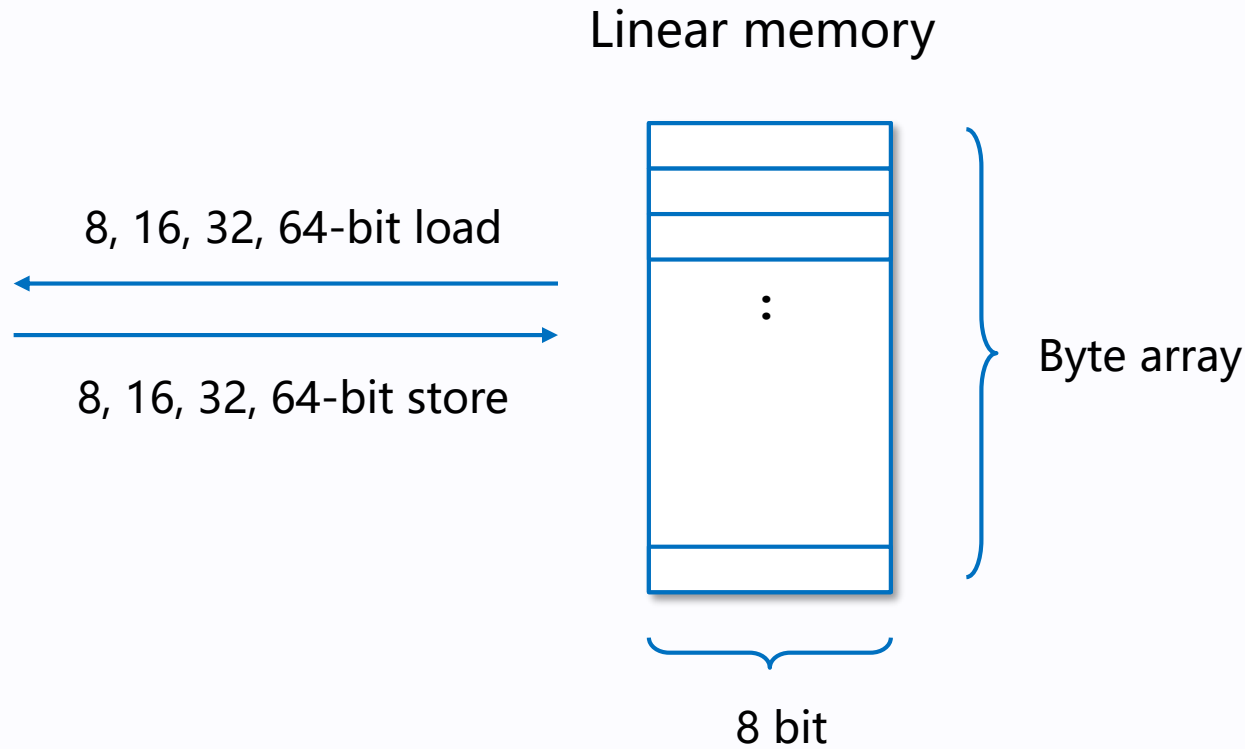
The function component of a module defines a vector of functions. Functions are referenced through function indices.

Table



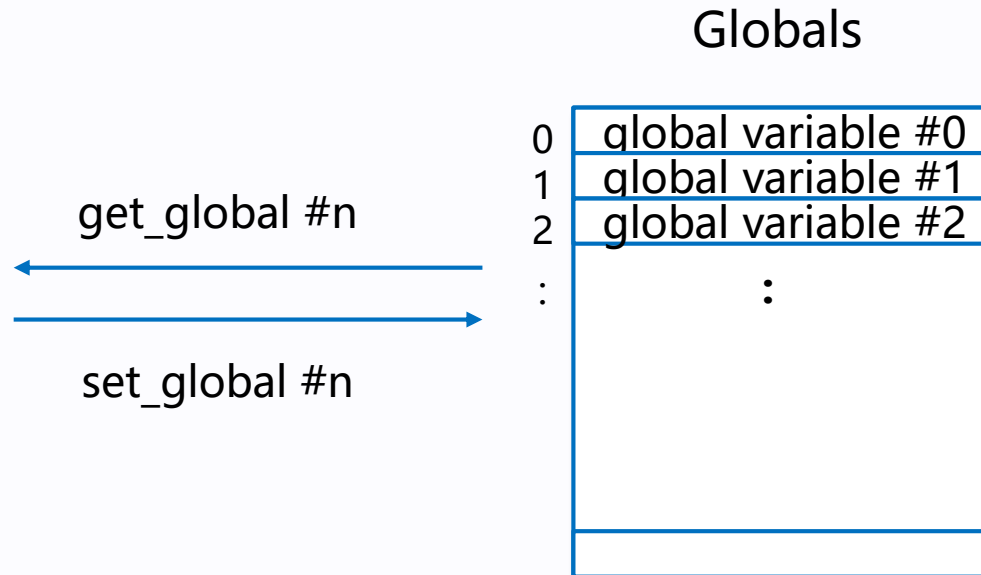
The table is an array of opaque values of a particular element type. Currently, the only available element type is an untyped function reference. This allows emulating function pointers by way of table indices. Tables are referenced through table indices.

Linear memory



The linear memory is a contiguous, mutable array of raw bytes.
The linear memory can be addressed at byte level (including unaligned).
The size of the memory is a multiple of the WebAssembly page size.

Globals



The globals component defines a vector of global variables.

The globals are referenced through global indices.

The global variables hold a value and can either be mutable or immutable.

2. WebAssembly abstract machine

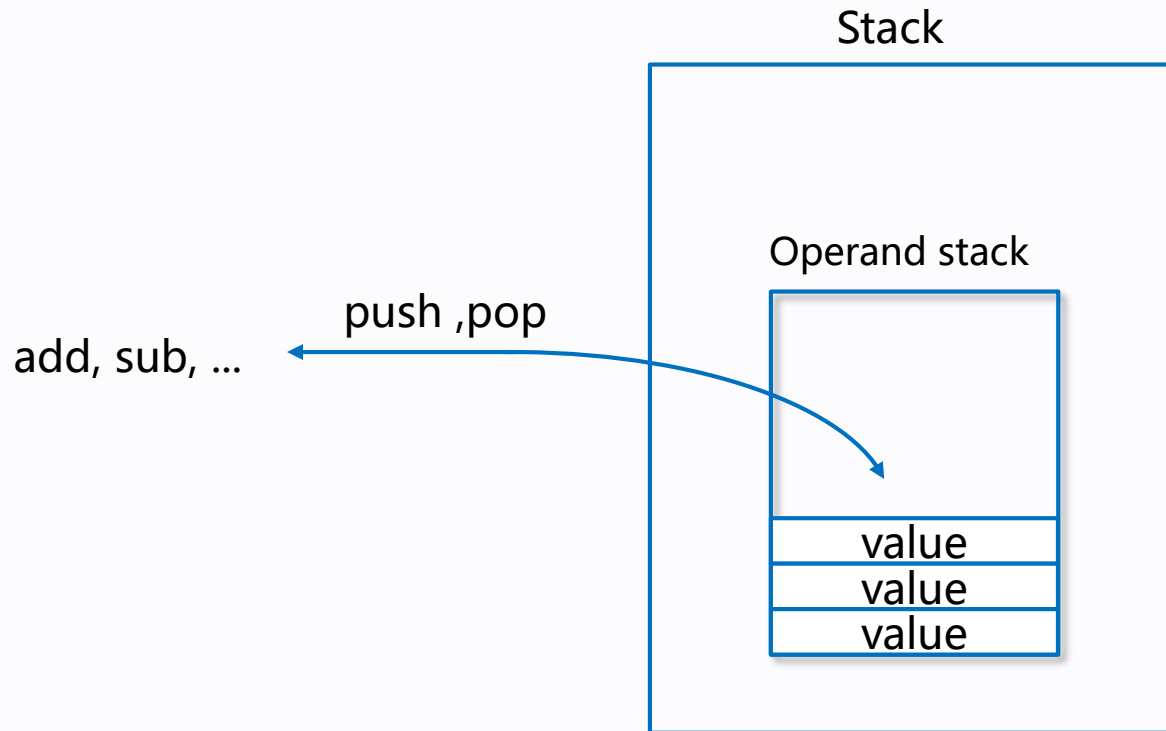
Stack

Stack



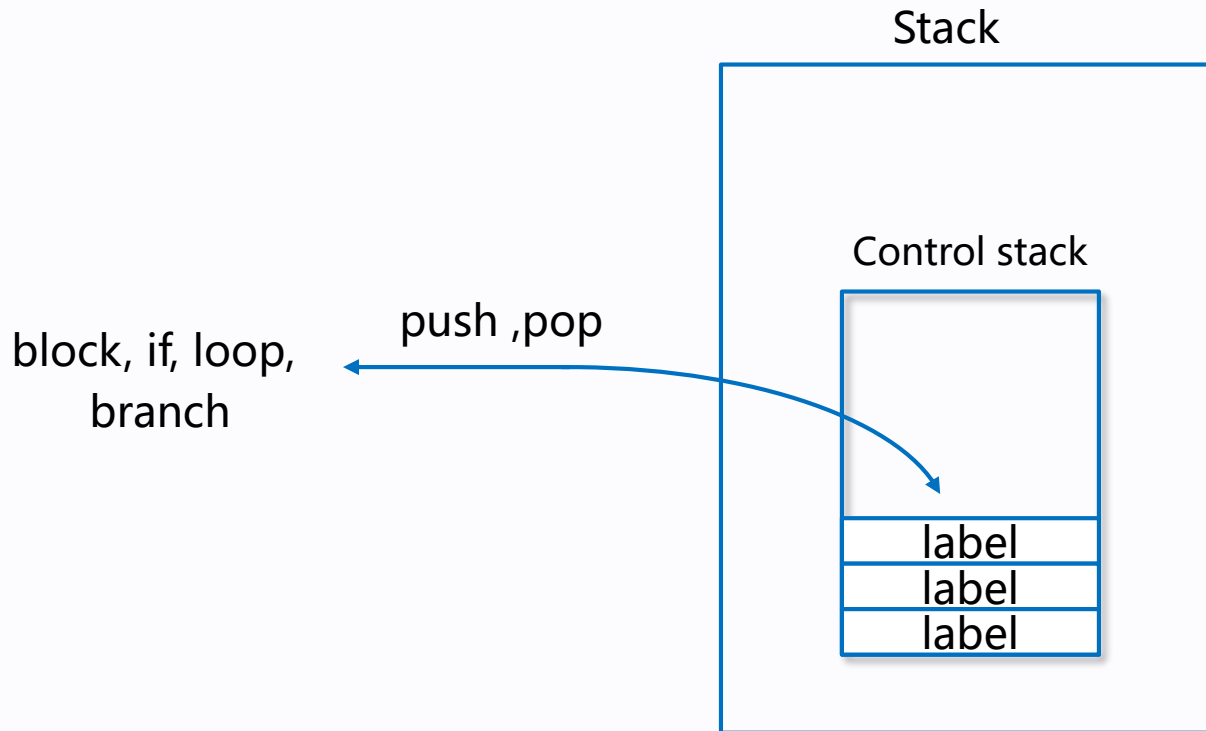
Most instructions interact with the implicit stack.
The stack contains values, labels and frames(activations).

Operand stack



Instructions manipulate values on an implicit operand stack.
The layout of the operand stack can be statically determined at any point in the code.

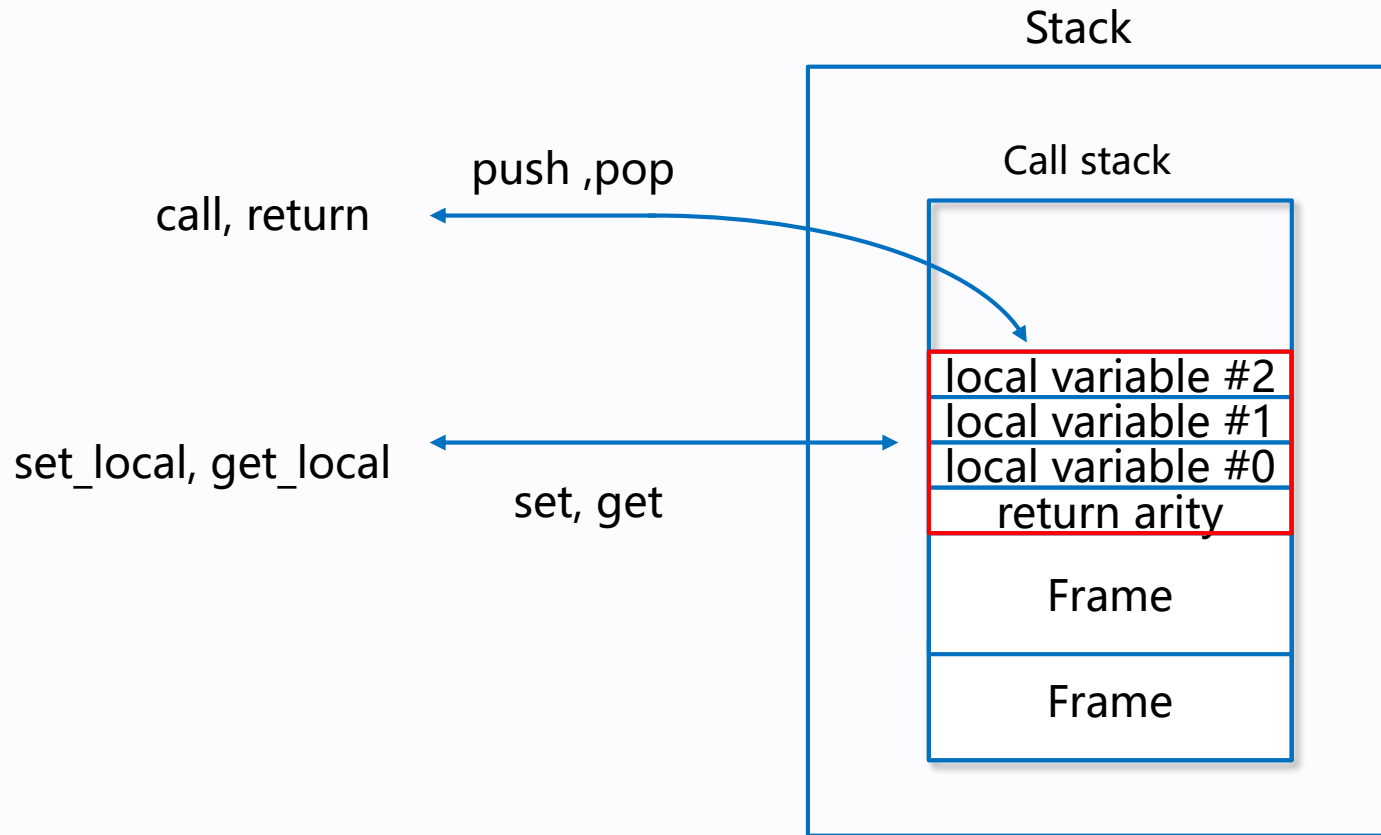
Control stack



Each structured control instruction introduces an implicit label.

Labels are targets for branch instructions that reference them with label indices.

Call stack



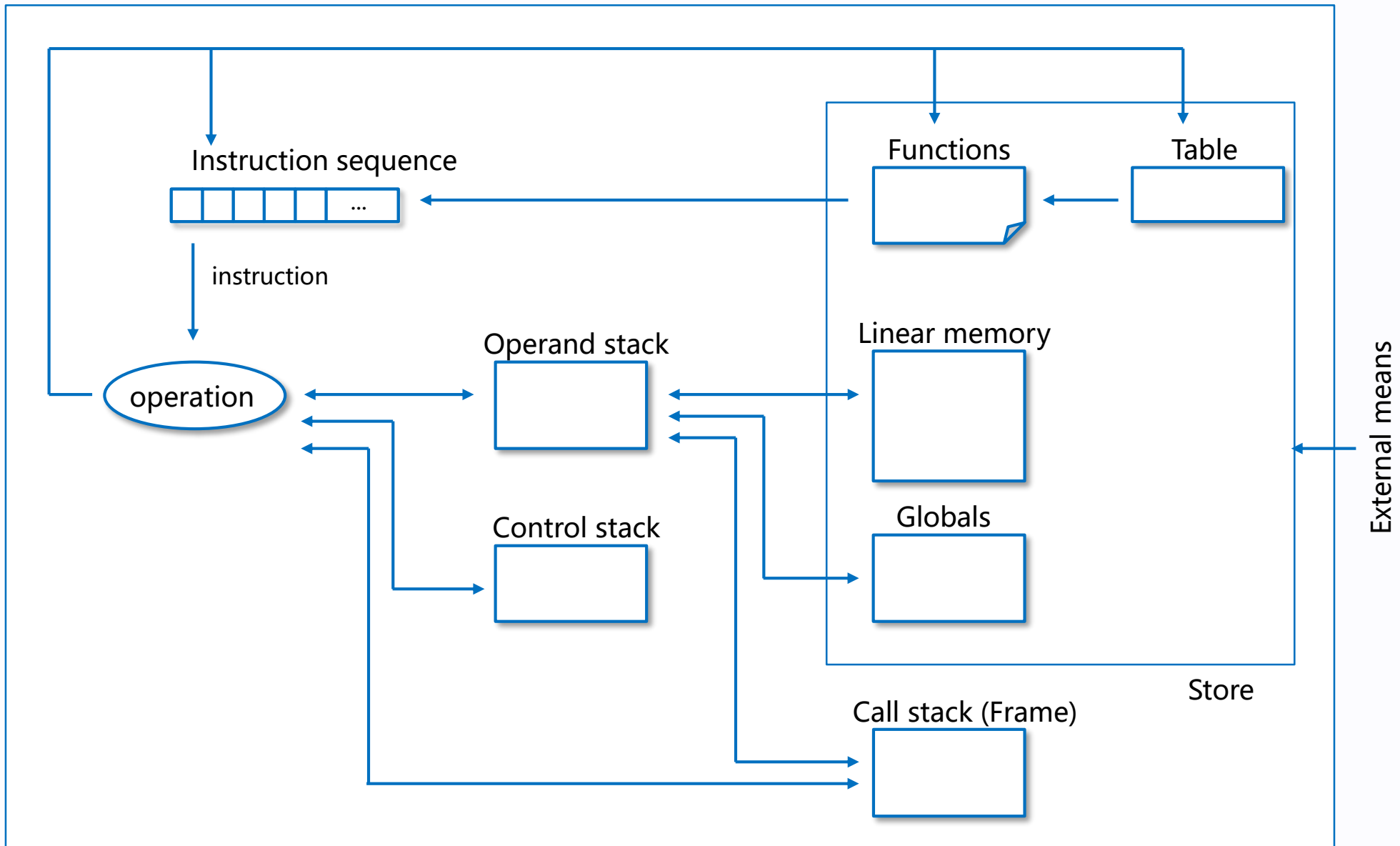
Frames hold the values of its local variables (including arguments).
Frames also carry the return arity of the respective function.

2. WebAssembly abstract machine

Computational model

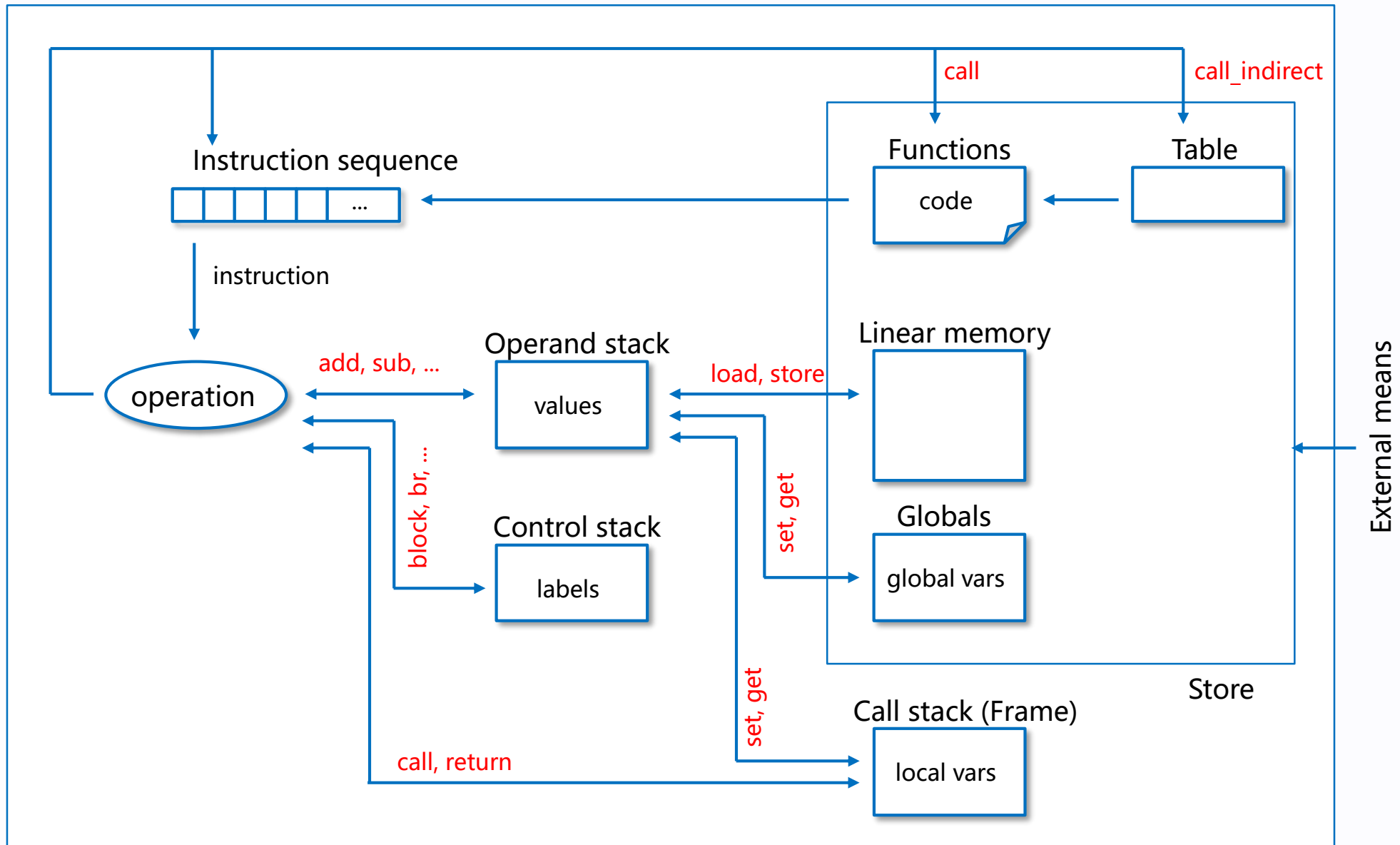
Computational model

WebAssembly abstract machine



Computational model

WebAssembly abstract machine



2. WebAssembly abstract machine

Type

Value types

Integers

32bit

Integers

64bit

Floating-point numbers

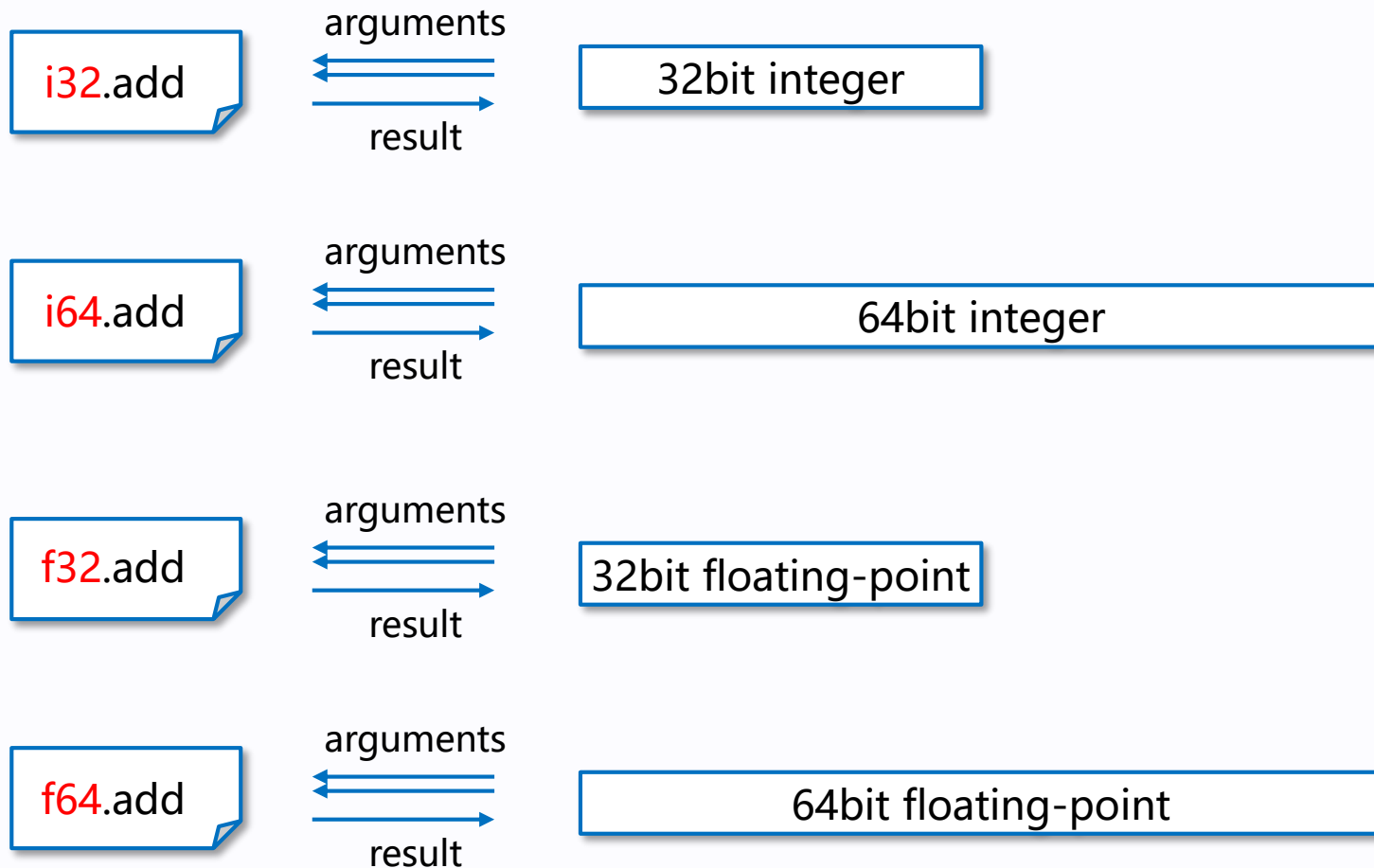
32bit

Floating-point numbers

64bit

WebAssembly provides only four basic value types.
32 bit integers also serve as Booleans and as memory addresses.

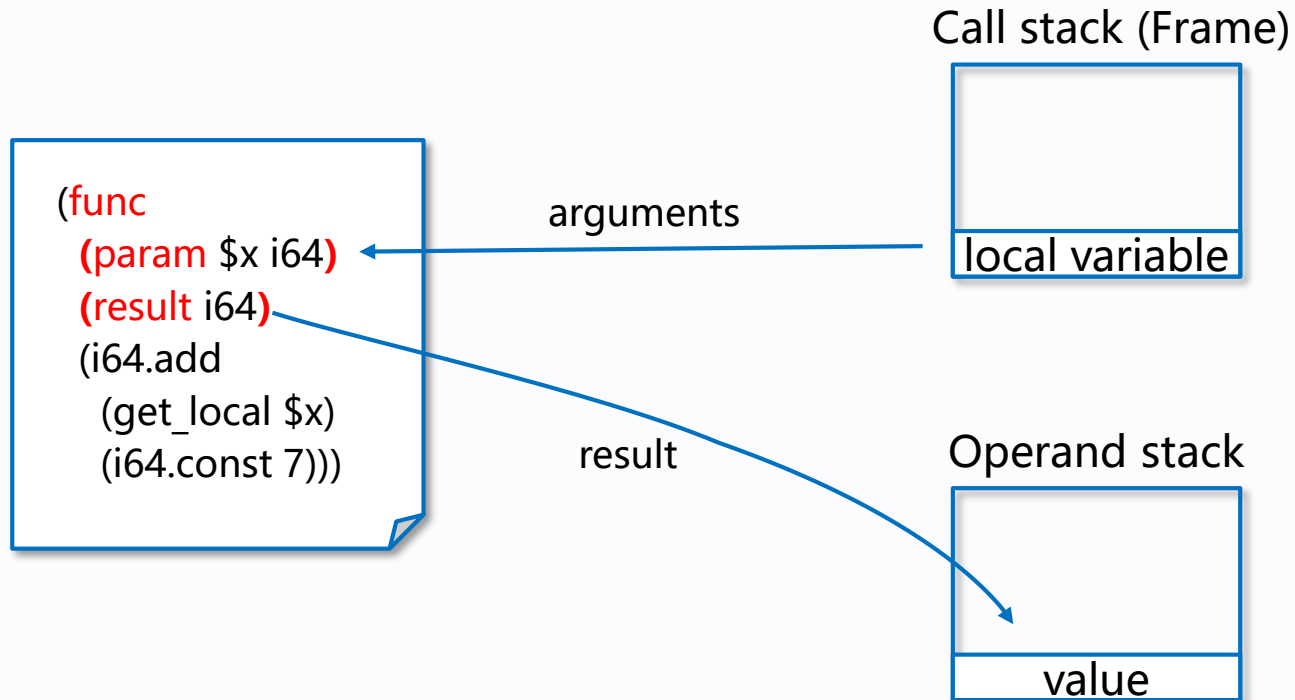
Instructions have type annotations



Some instructions have type annotations.

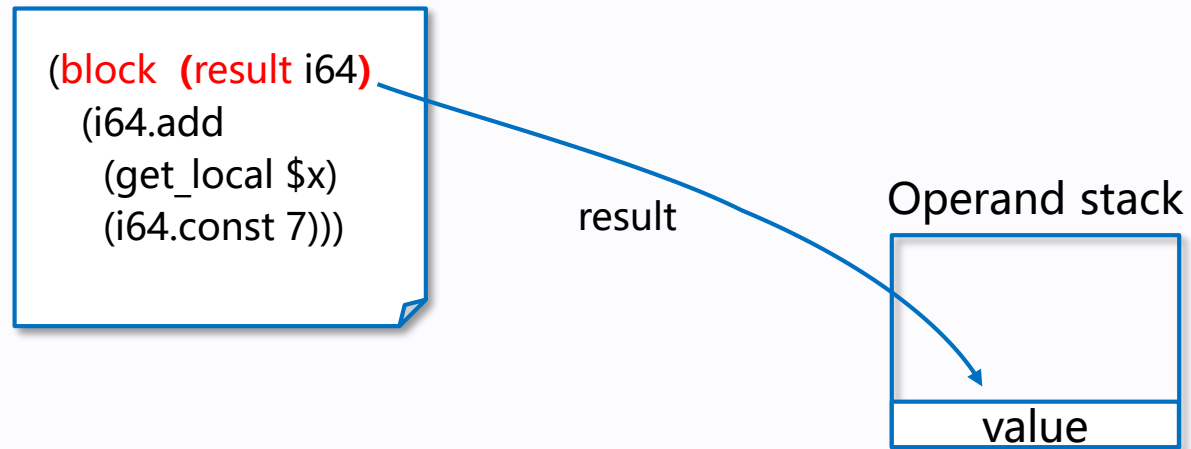
For example, the instruction `i32.add` has type `[i32 i32] → [i32]`, consuming two `i32` values and producing one.

Functions have type declarations



Each function takes a sequence of WebAssembly values as parameters and returns a sequence of values as results as defined by its function type.

Control blocks have also a type declaration



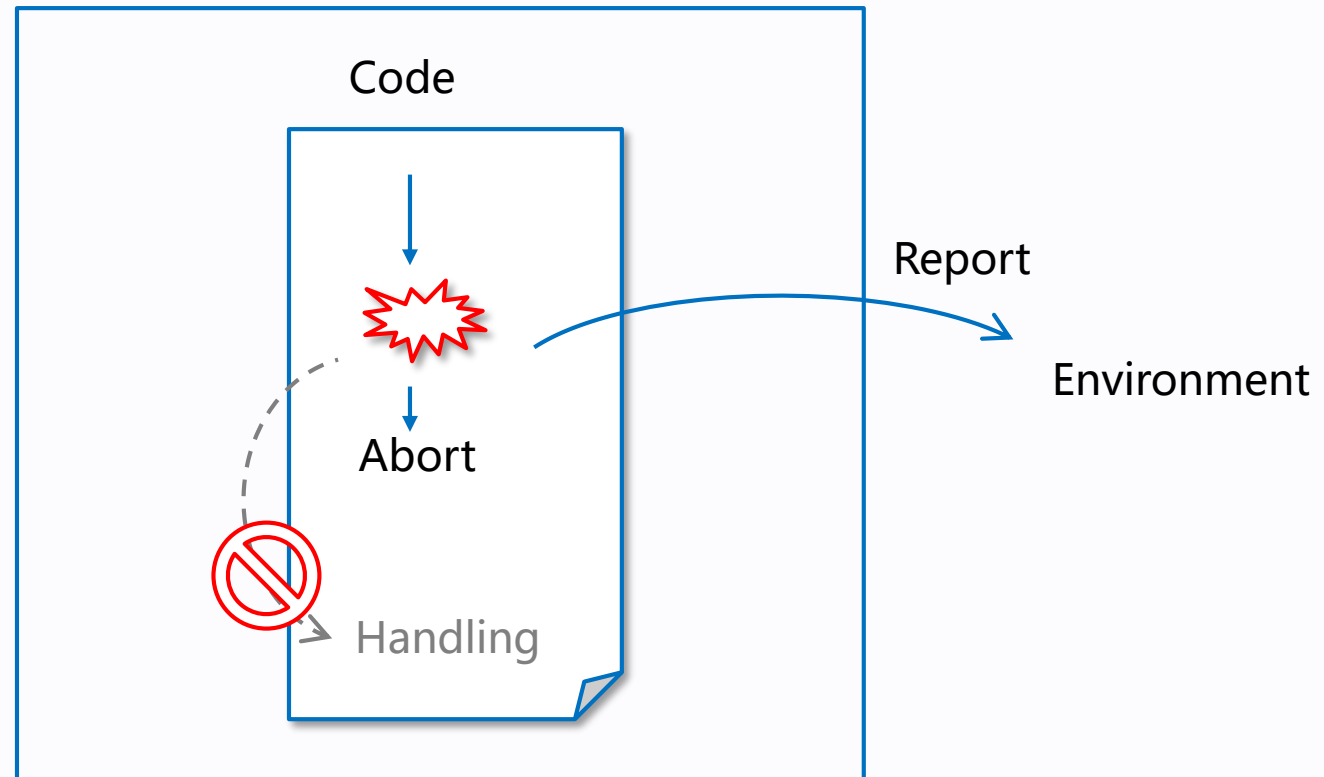
Every control construct is annotated with a function type.

2. WebAssembly abstract machine

Trap

Trap

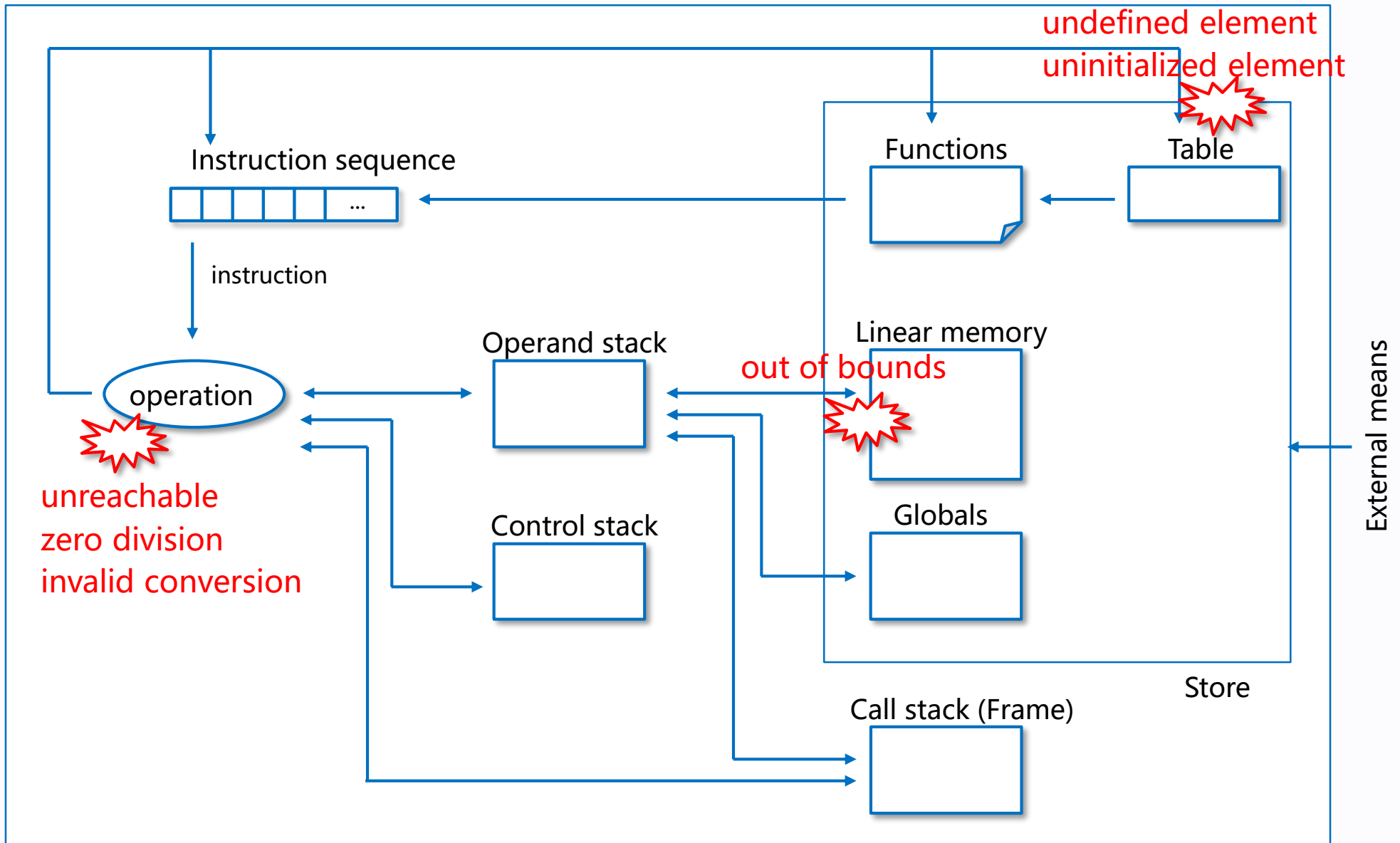
WebAssembly abstract machine



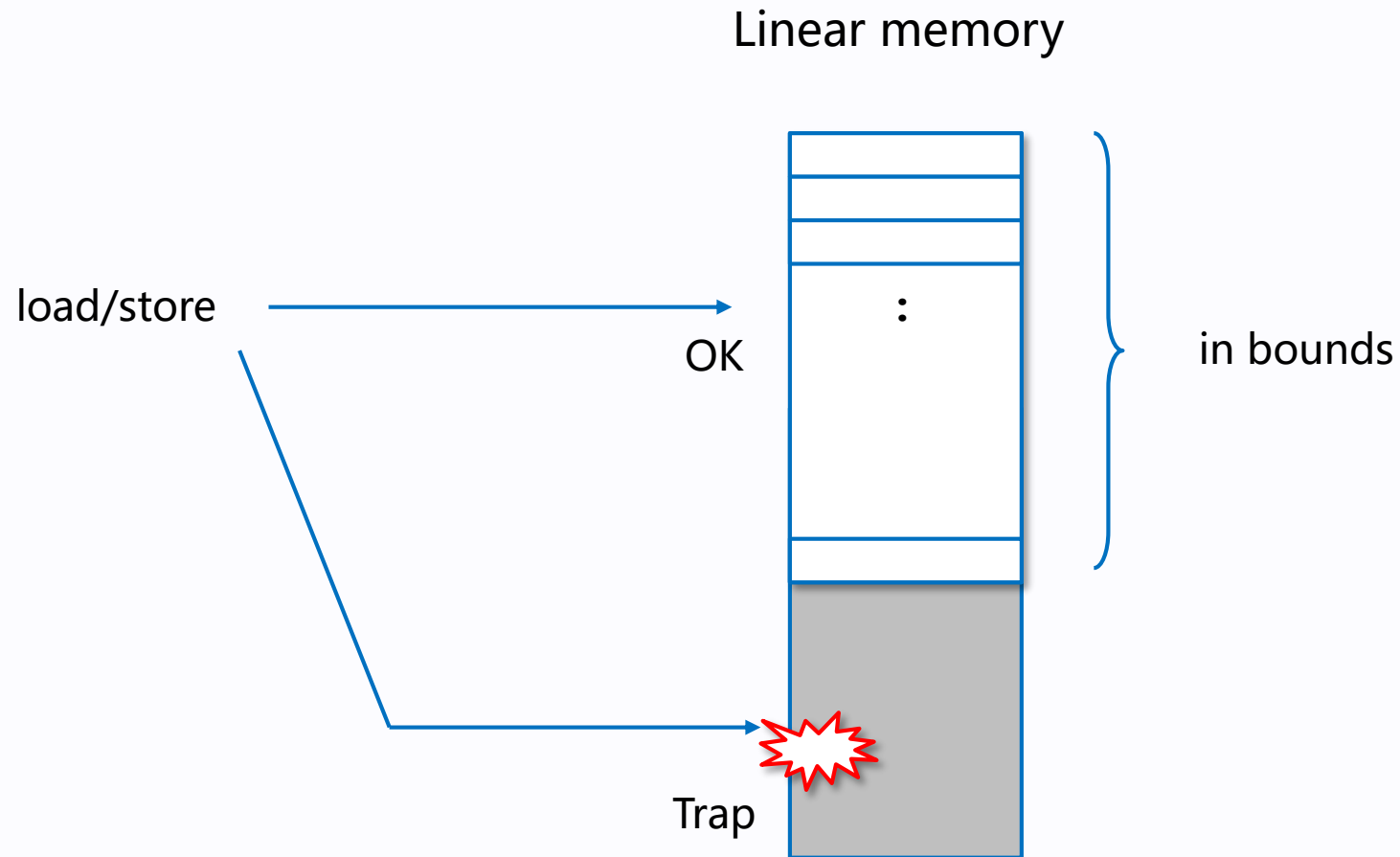
Certain instructions may produce a trap, which immediately aborts execution. Traps cannot be handled by WebAssembly code, but are reported to the outside environment, where they typically can be caught.

Trap

WebAssembly abstract machine



Linear memory



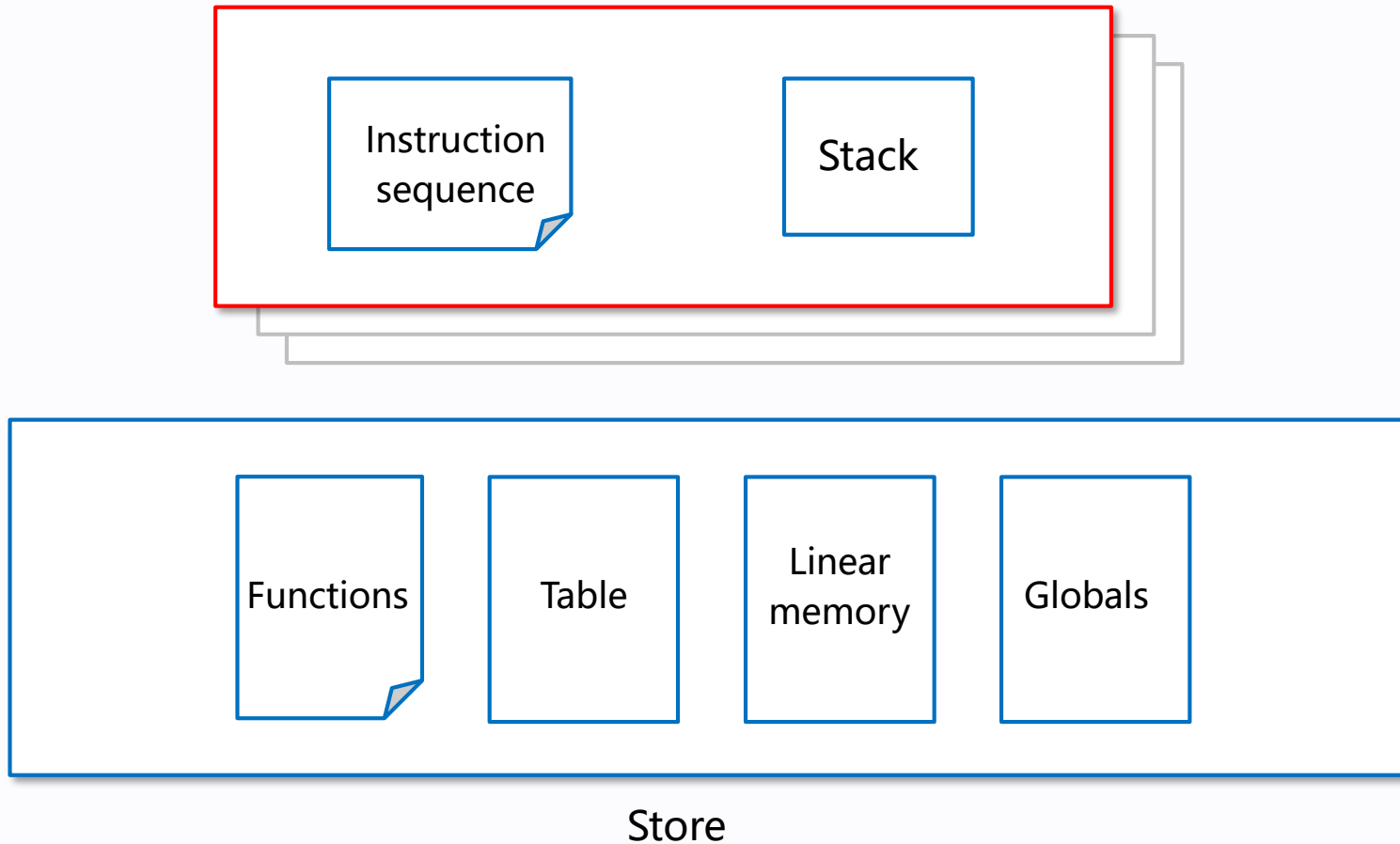
A trap occurs if an access is not within the bounds of the current memory size.

2. WebAssembly abstract machine

Thread

Thread

Threads



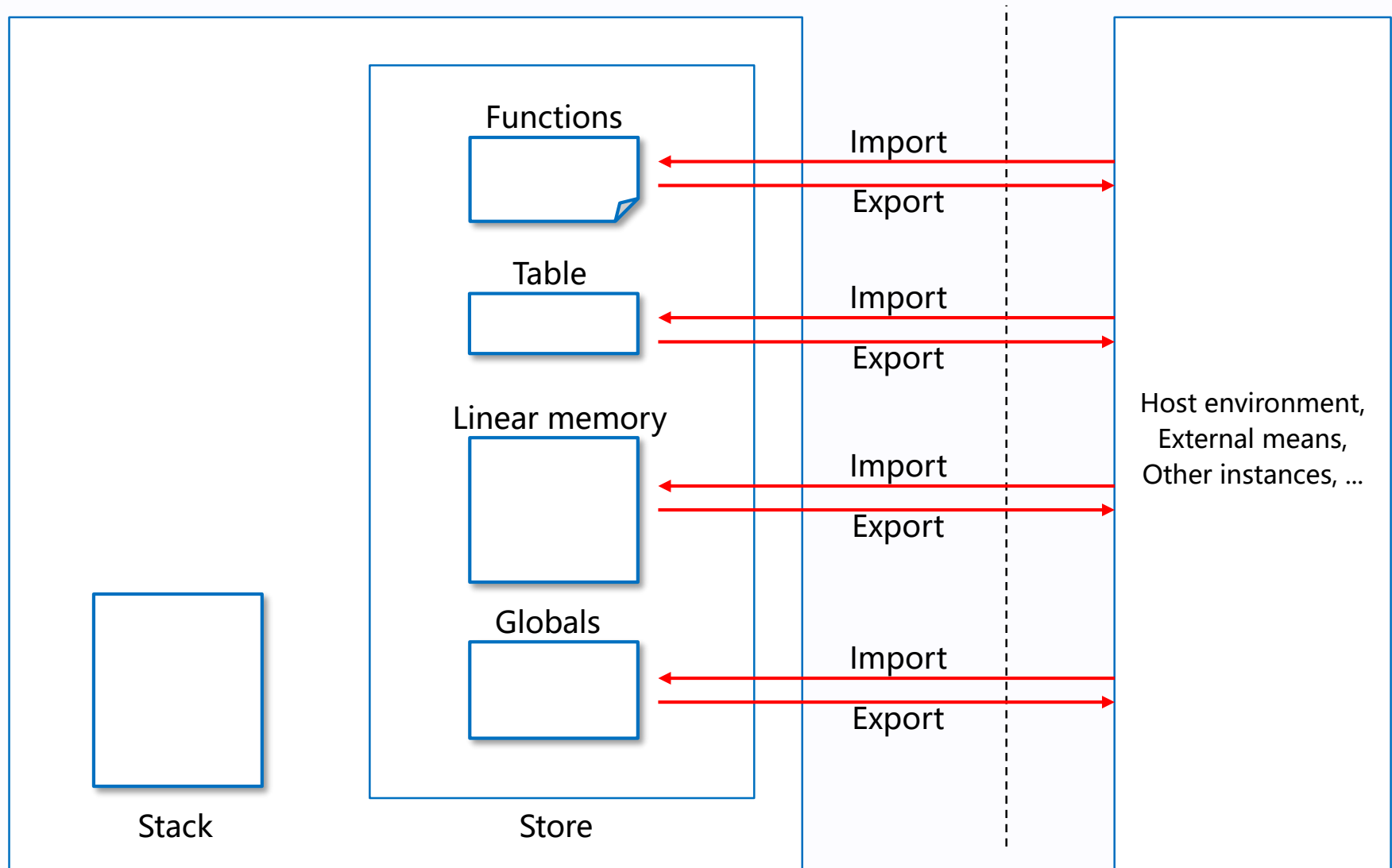
The current version of WebAssembly is single-threaded, but configurations with multiple threads may be supported in the future.

2. WebAssembly abstract machine

External interface

Import and export

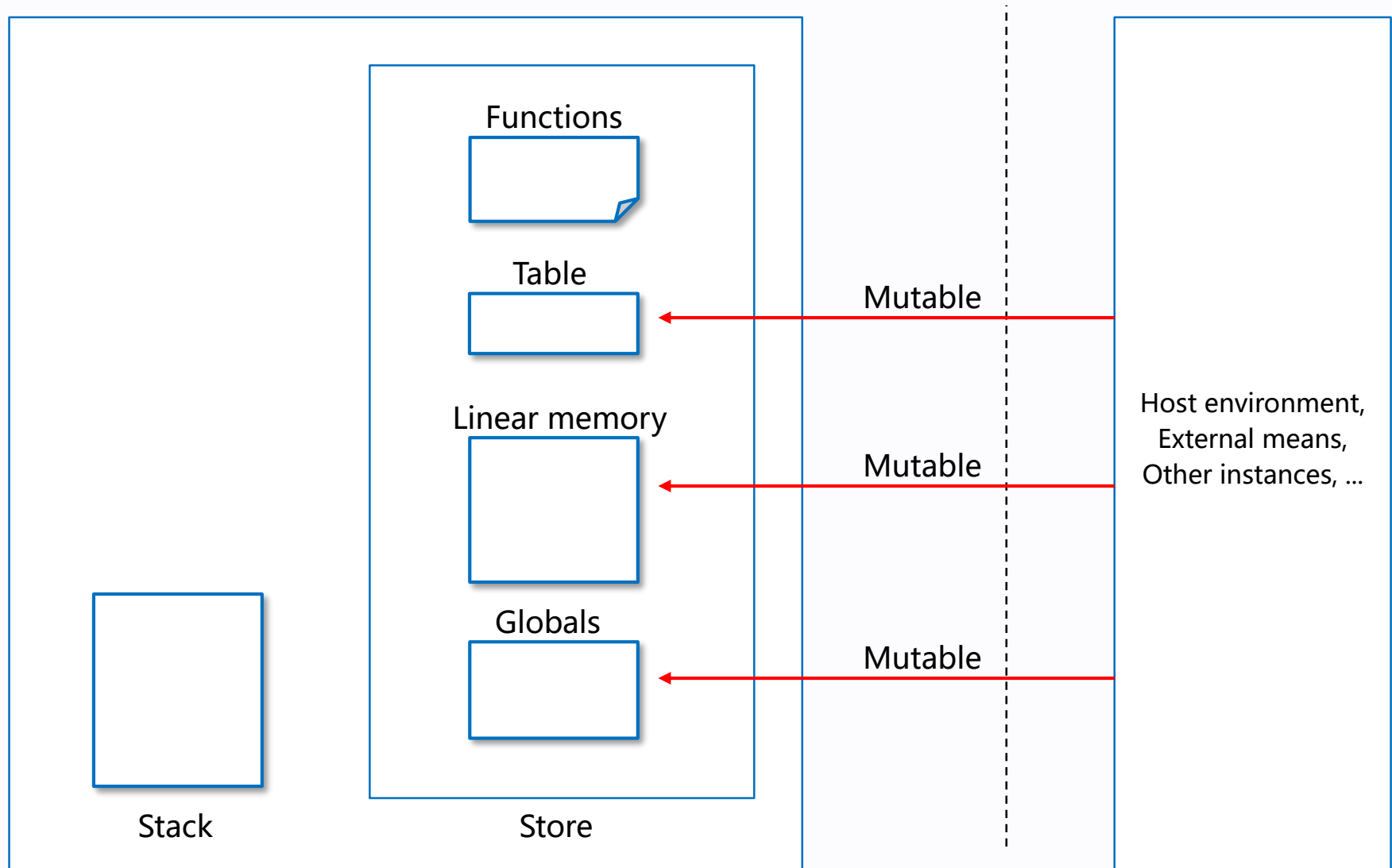
WebAssembly abstract machine



Functions, table, memory and globals may be shared via import/export.

Mutation from external

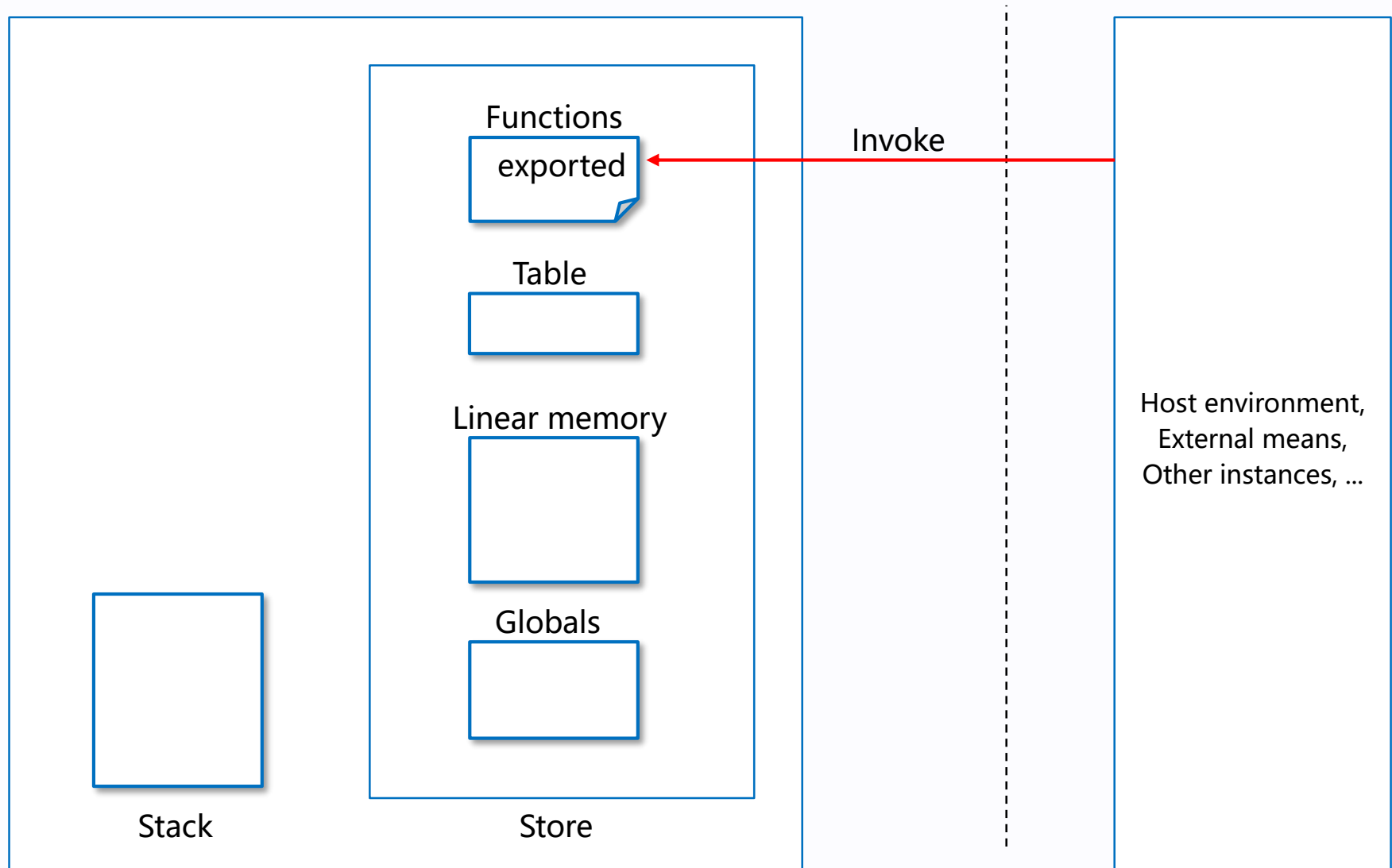
WebAssembly abstract machine



Table, memory and globals can be mutated by external mean.

Invoke from external

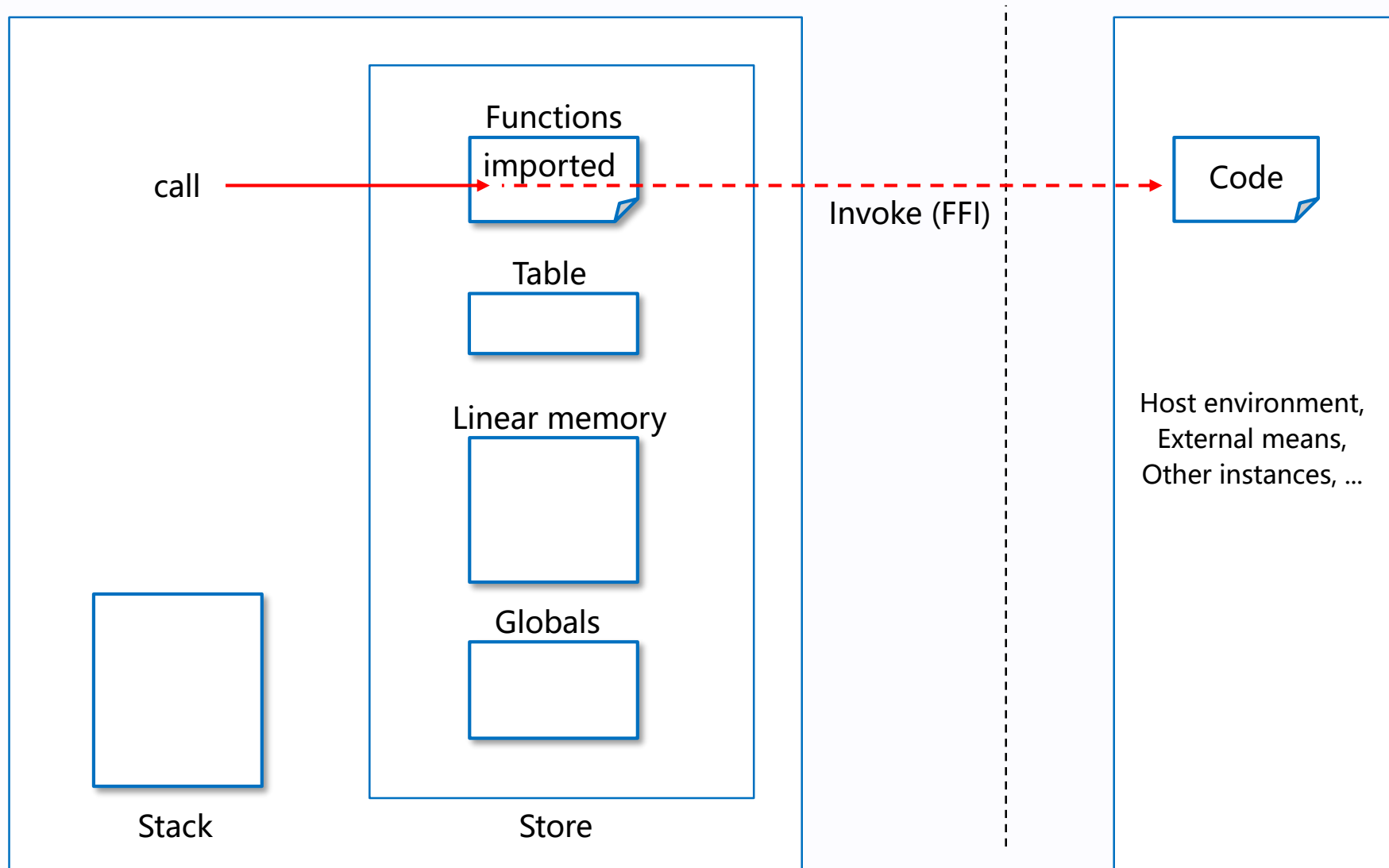
WebAssembly abstract machine



Any exported function can be invoked externally.

Foreign call

WebAssembly abstract machine



Call instructions can invoke an imported function.

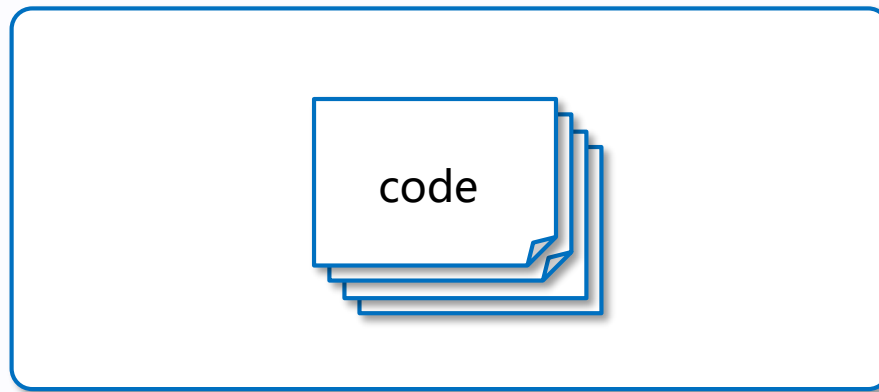
3. WebAssembly module

3. WebAssembly module

Module

WebAssembly module

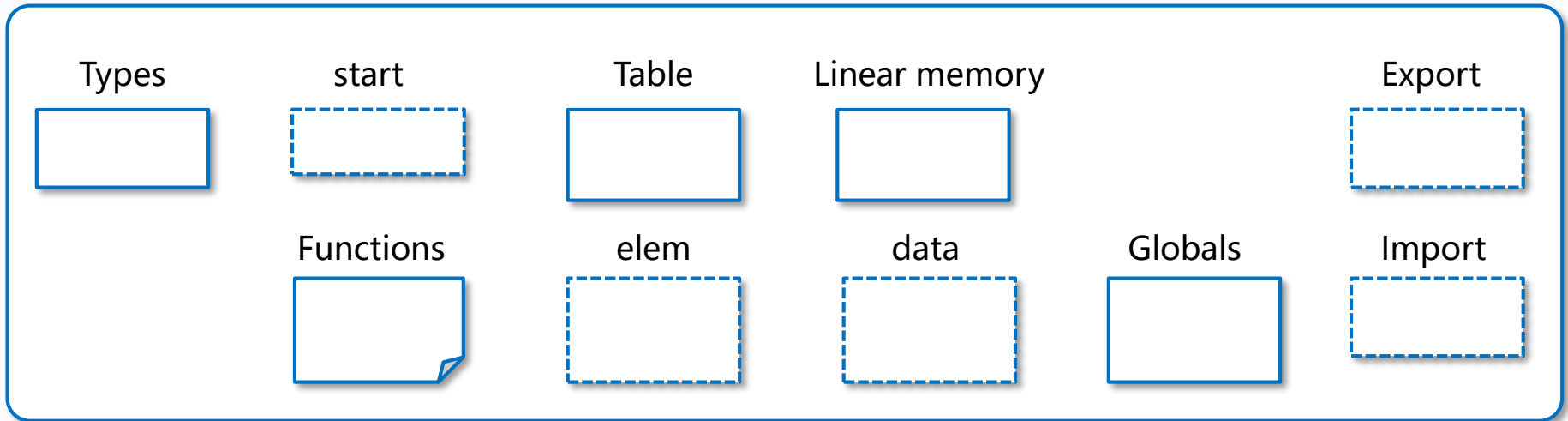
WebAssembly module



WebAssembly programs are organized into modules.
Modules are the distributable, loadable, and executable unit of code.
WebAssembly modules are distributed in a binary format.

WebAssembly module

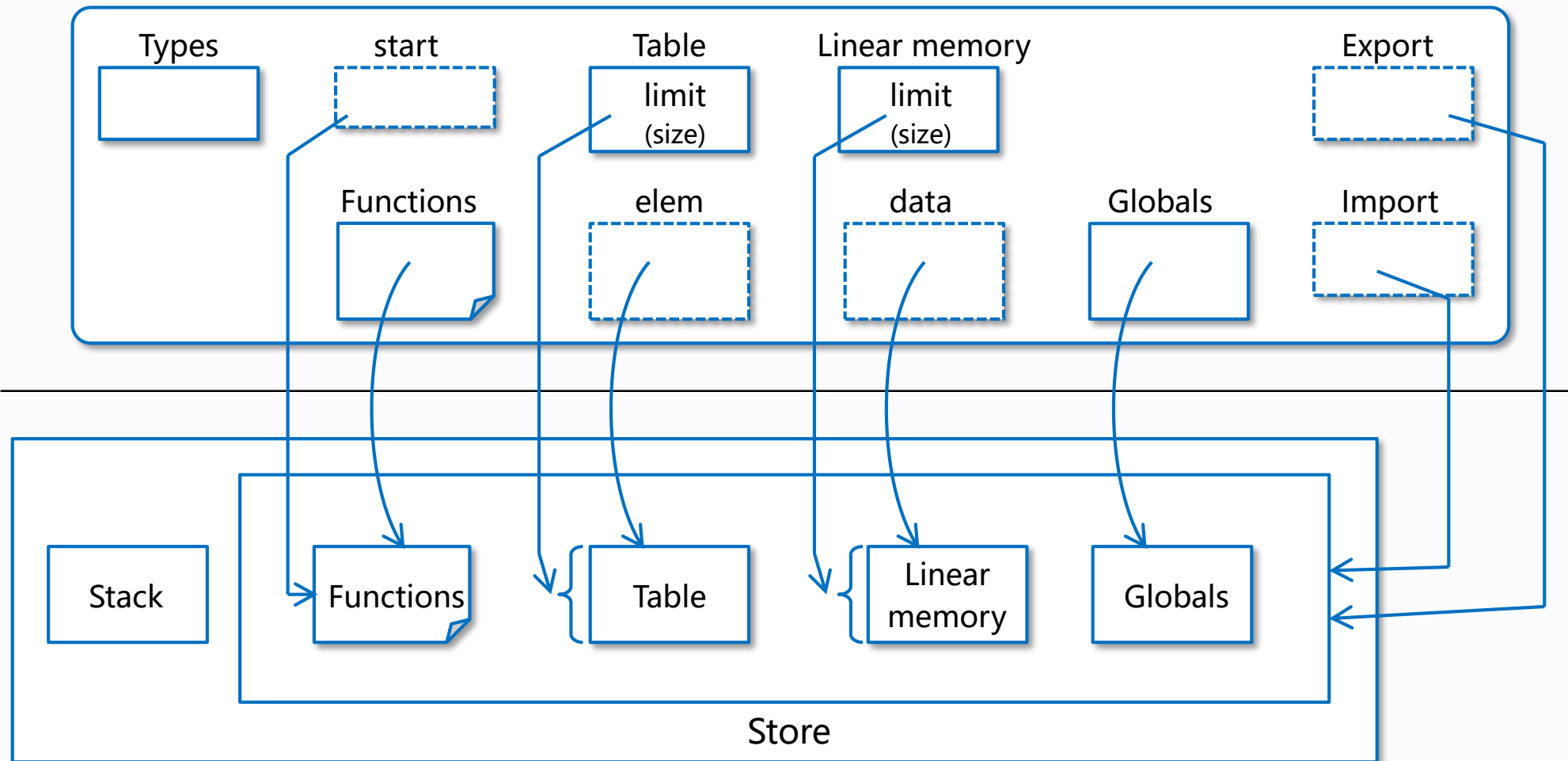
WebAssembly module



A module collects definitions for types, functions, table, memory, and globals. In addition, it can declare imports and exports and provide initialization logic in the form of data and element segments or a start function.

WebAssembly module and abstract machine

WebAssembly **module**



WebAssembly **abstract machine** (module instance)

A module corresponds to the static representation of a program.
A module instance corresponds to a dynamic representation.

3. WebAssembly module

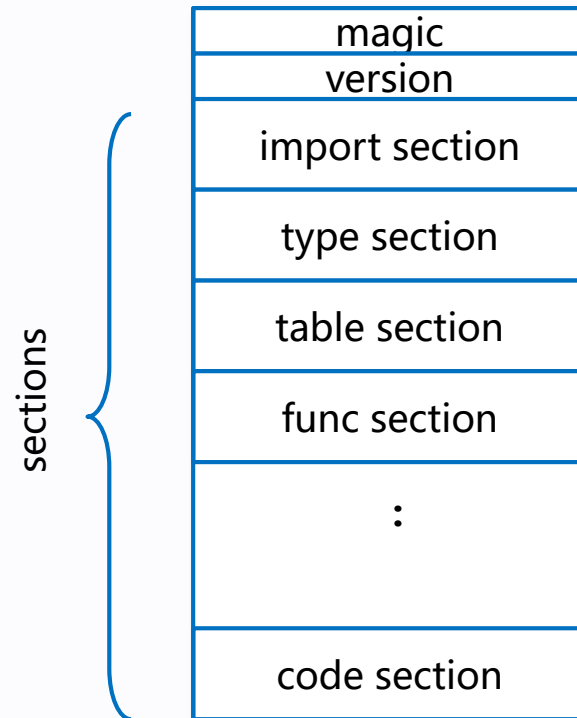
Binary encoding

Binary encoding of modules

WebAssembly module (WebAssembly binary)

00 61 73 6d 01 00 00 00 01 07 01 60 02 7e 7e 01 7e 03 ...

Form

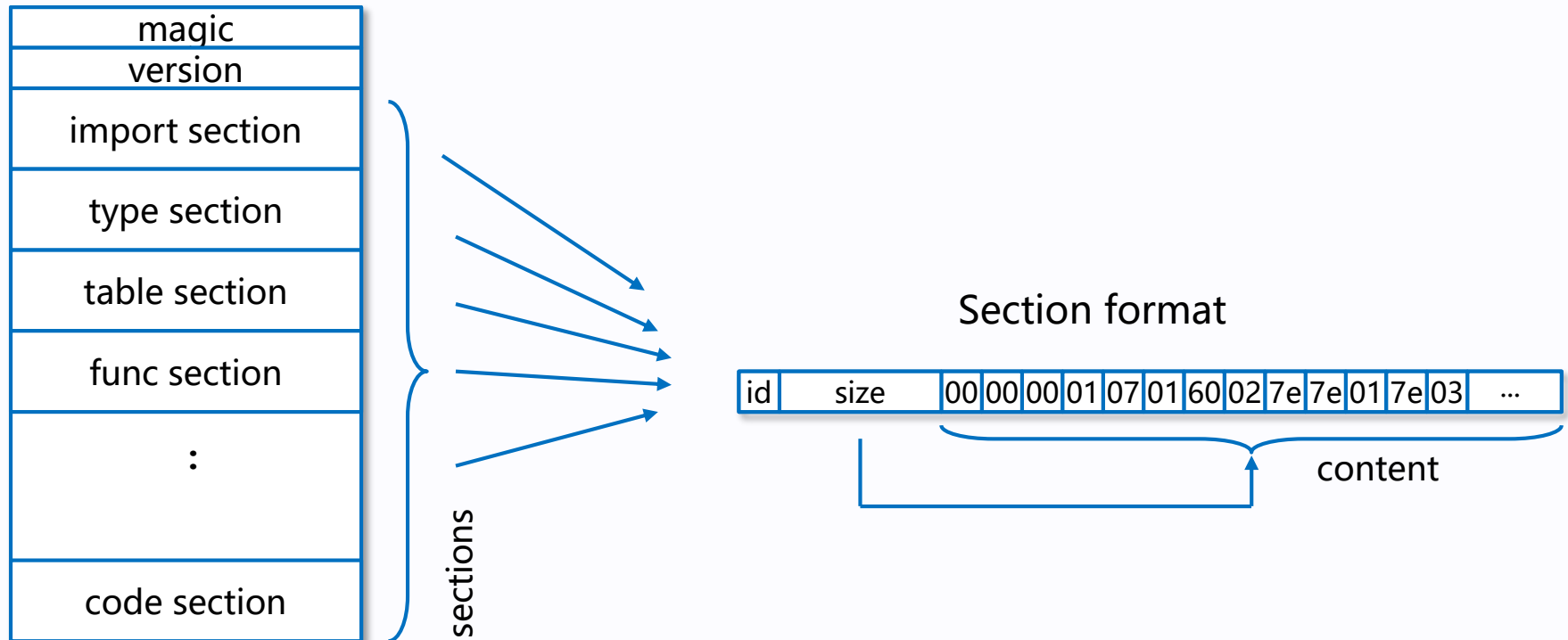


Binary encoding of modules

The binary encoding of modules is organized into sections.

Sections

Binary encoding of modules




Each section consists of

- a one-byte section id,
- the u32 size of the contents, in bytes,
- the actual contents, whose structure is depended on the section id.

Example of WebAssembly module

[text format]

```
(module
  (func (export "foo" (result i32)
    i32.const 7))
```

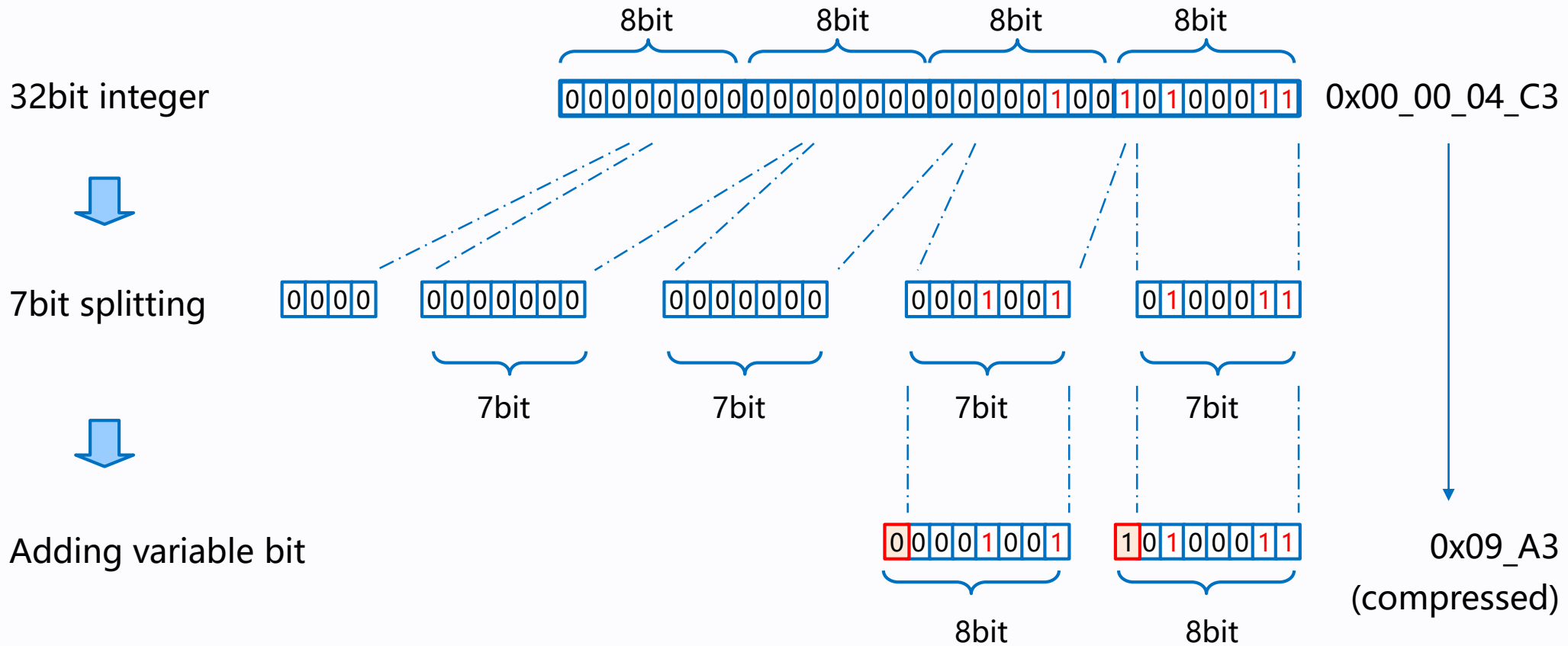
 (wat2wasm -v` command)

[binary format]

```
0000000: 0061 736d      ; WASM_BINARY_MAGIC
0000004: 0100 0000      ; WASM_BINARY_VERSION
; section "Type" (1)
0000008: 01              ; section code
0000009: 05              ; section size
000000a: 01              ; num types
; type 0
000000b: 60              ; func
000000c: 00              ; num params
000000d: 01              ; num results
000000e: 7f              ; i32
; section "Function" (3)
000000f: 03              ; section code
0000010: 02              ; section size
0000011: 01              ; num functions
0000012: 00              ; function 0 signature
; index
```

```
; section "Export" (7)
0000013: 07              ; section code
0000014: 07              ; section size
0000015: 01              ; num exports
0000016: 03              ; string length
0000017: 666f 6f        ; foo ; export name
000001a: 00              ; export kind
000001b: 00              ; export func index
; section "Code" (10)
000001c: 0a              ; section code
000001d: 06              ; section size
000001e: 01              ; num functions
; function body 0
000001f: 04              ; func body size
0000020: 00              ; local decl count
0000021: 41              ; i32.const
0000022: 07              ; i32 literal
0000023: 0b              ; end
```

Integer encoding with LEB128



All integers are encoded using the LEB128 variable-length integer encoding.

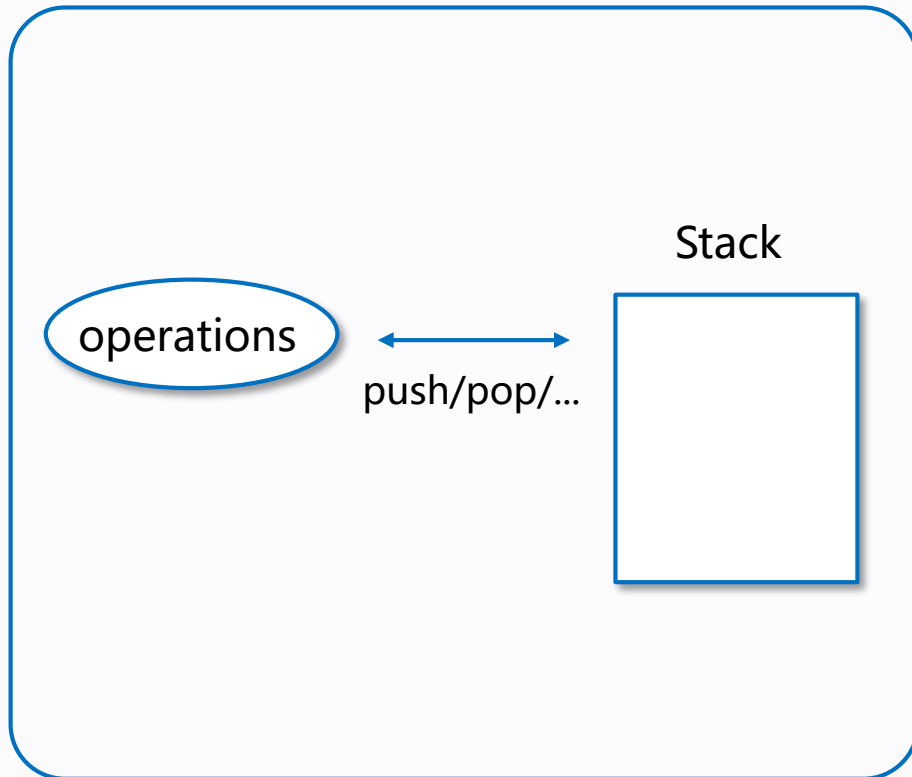
4. WebAssembly instructions

4. WebAssembly instructions

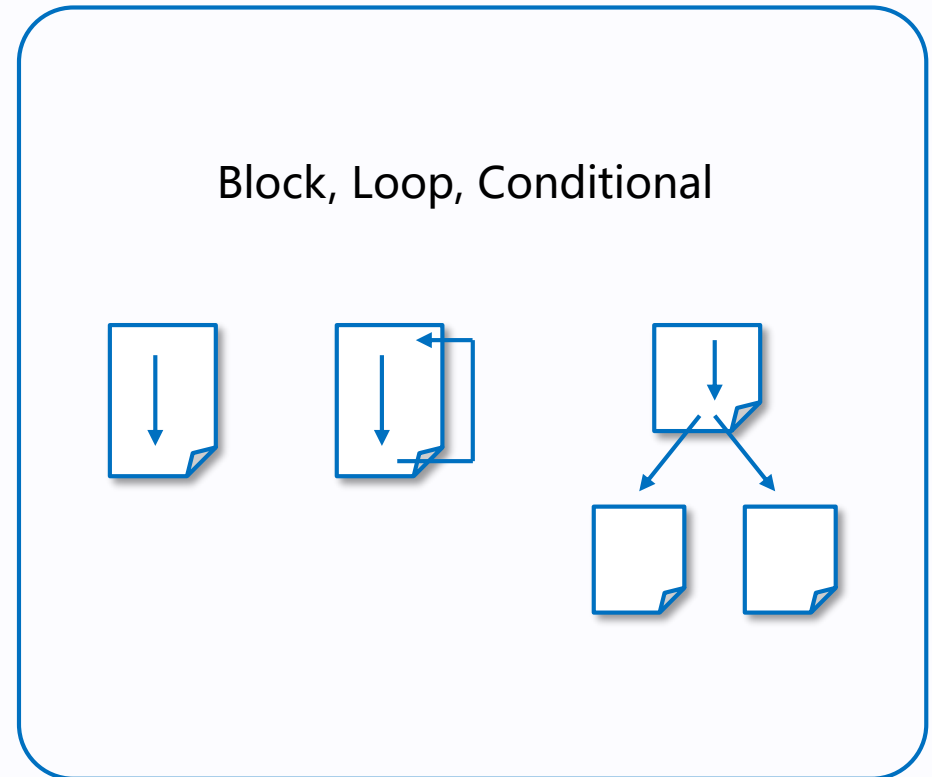
Instructions

Instructions

Simple instructions



Control instructions

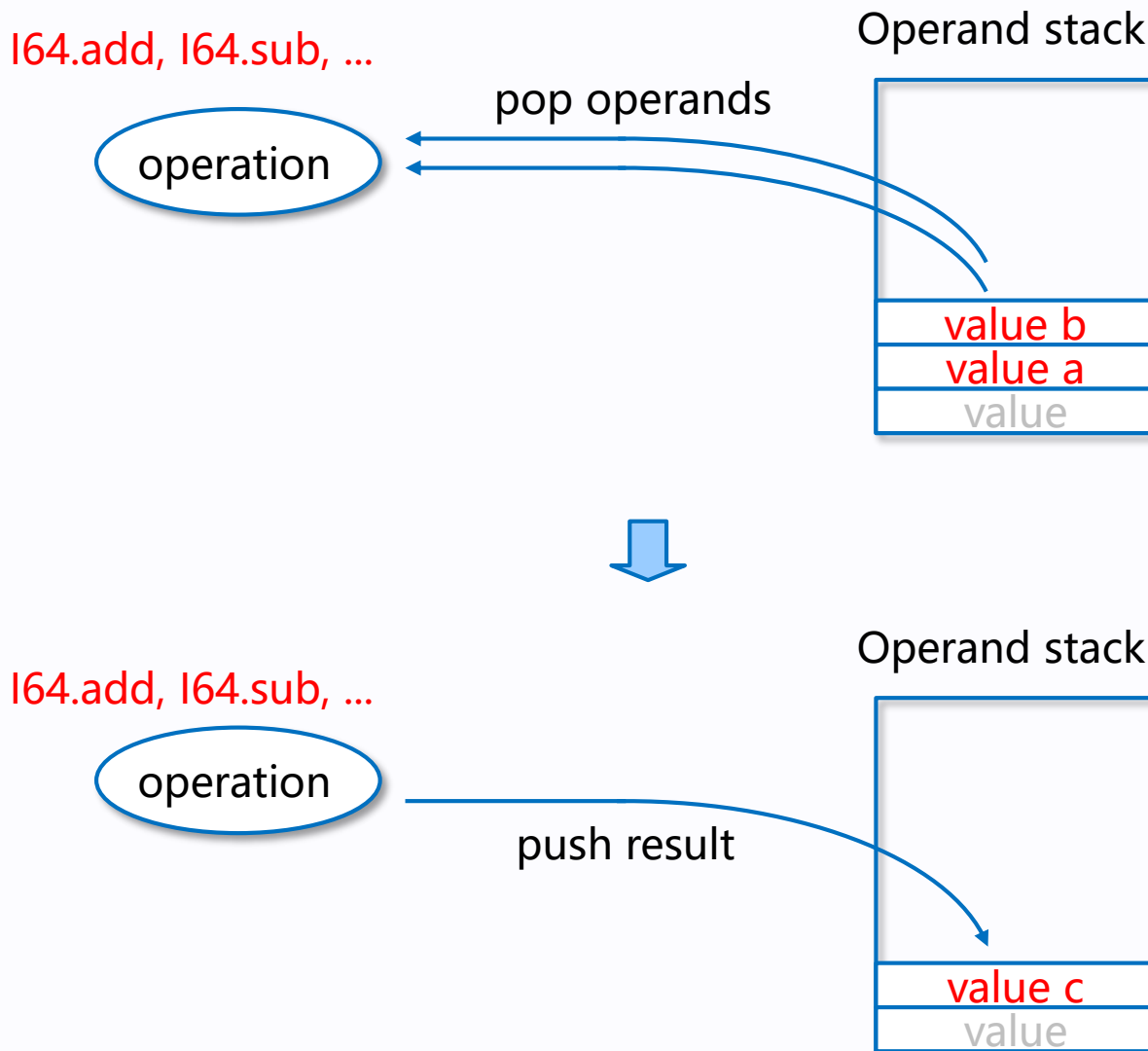


Instructions fall into two main categories.
Simple instructions perform basic operations on data.
Control instructions alter control flow.

4. WebAssembly instructions

Simple instructions

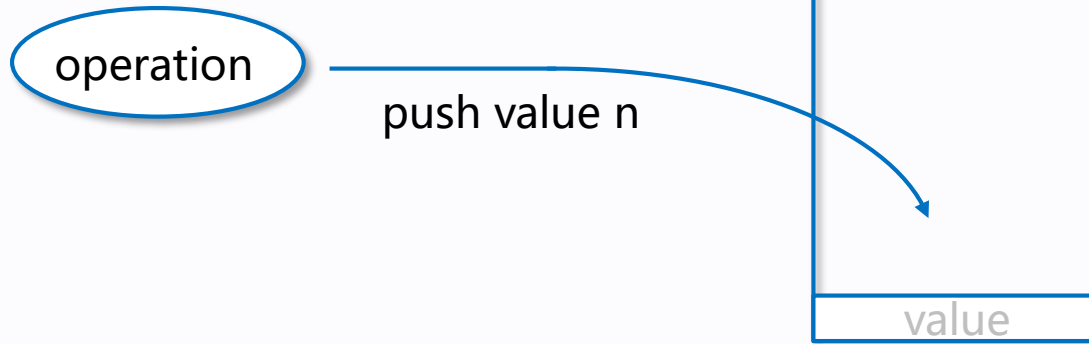
Numeric instructions



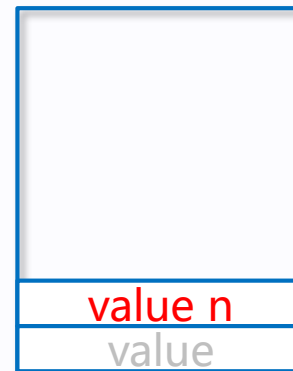
Numeric instructions pop arguments from the operand stack and push results back to it.

Numeric instructions : const

l64.const n, l32.const n, ...

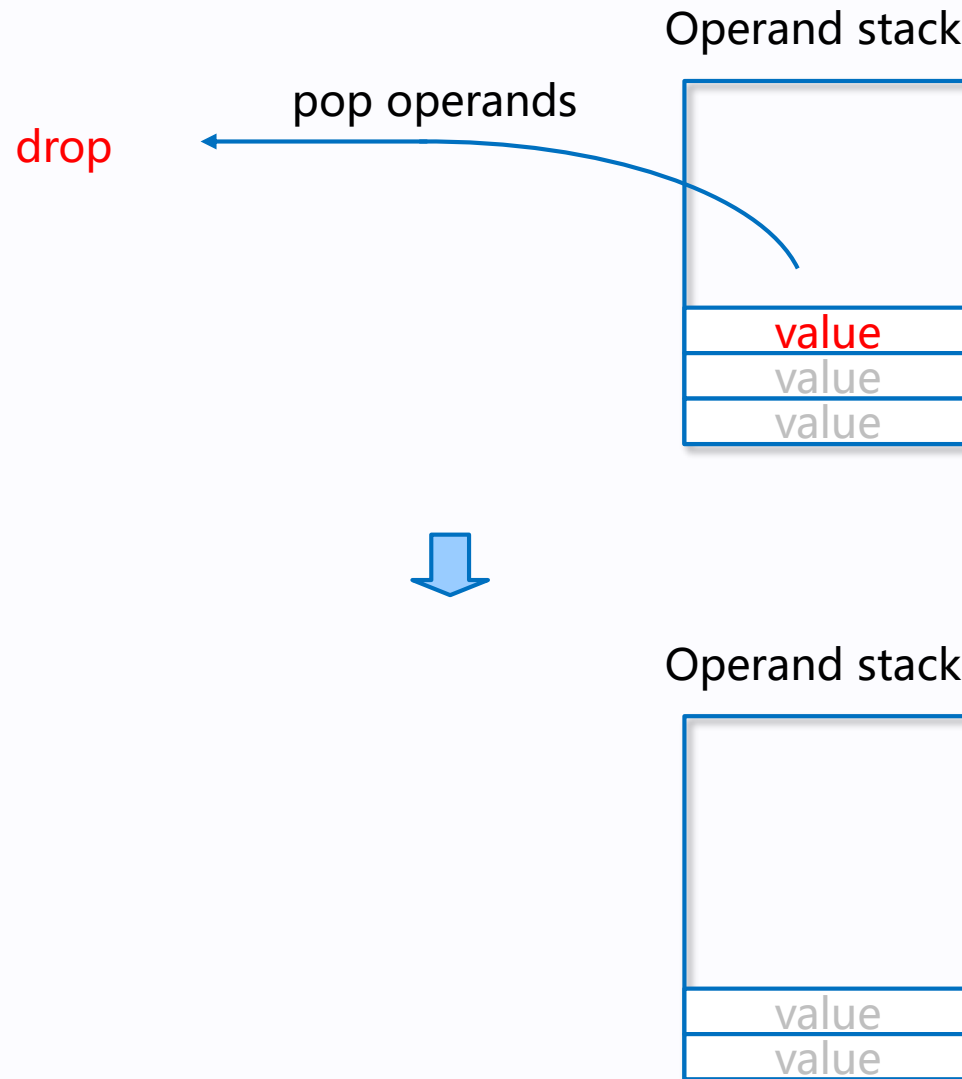


Operand stack



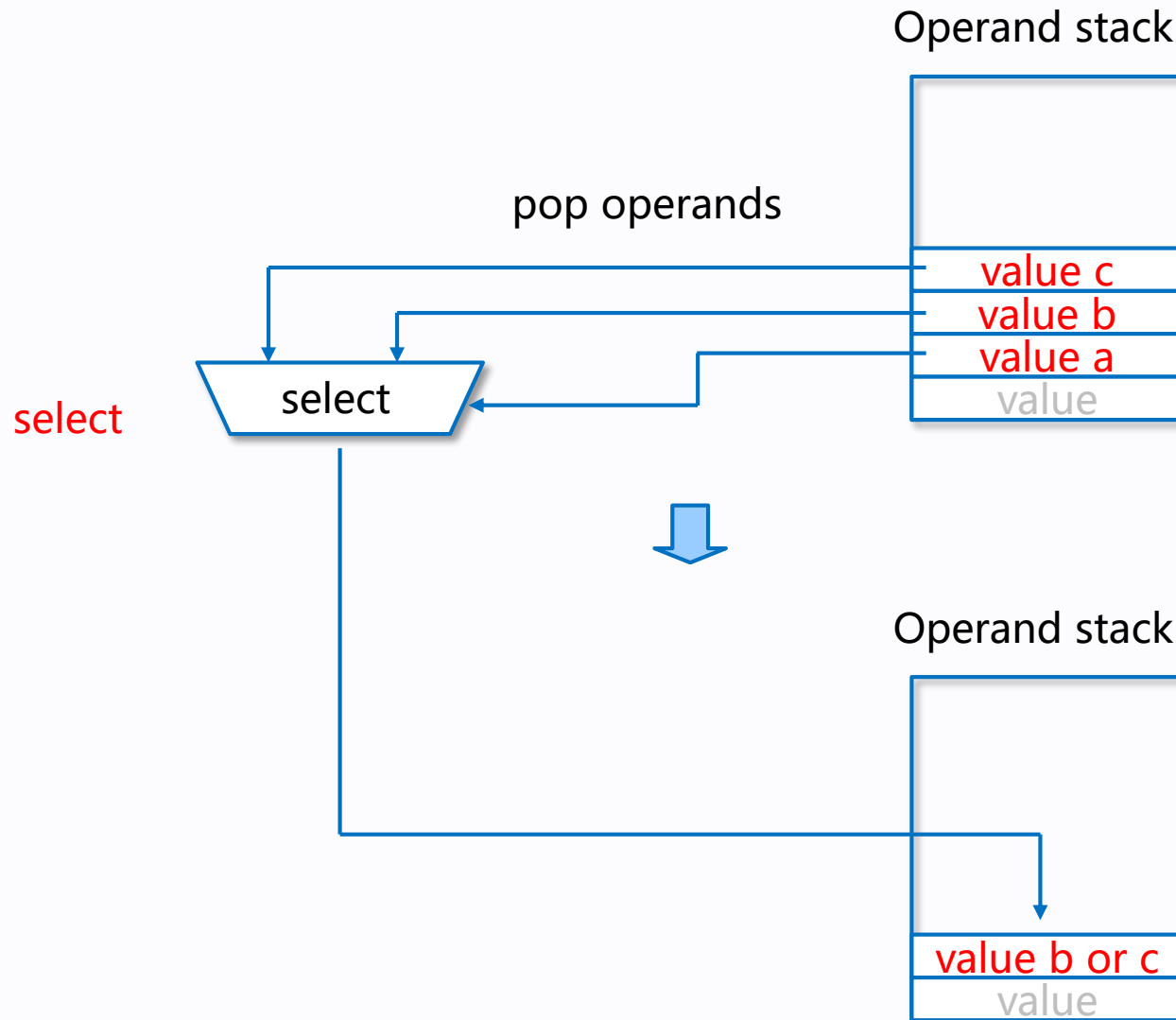
The const instruction pushes the value to the stack.

Parametric instructions : drop



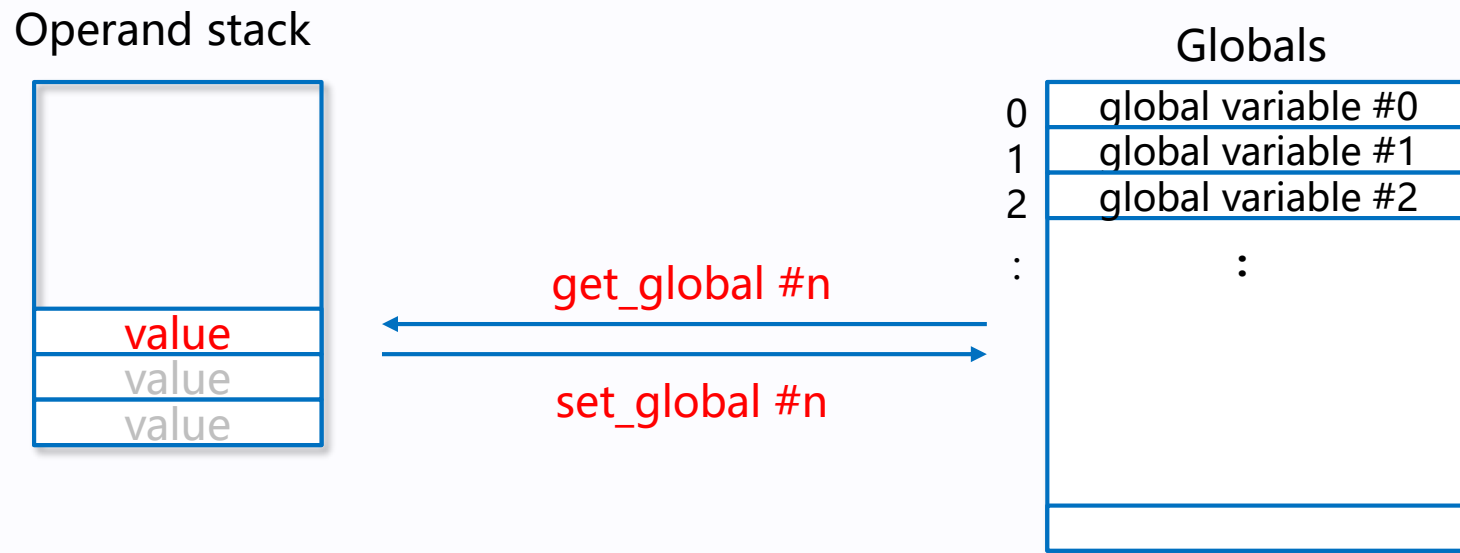
The drop instruction simply throws away a single operand.

Parametric instructions : select



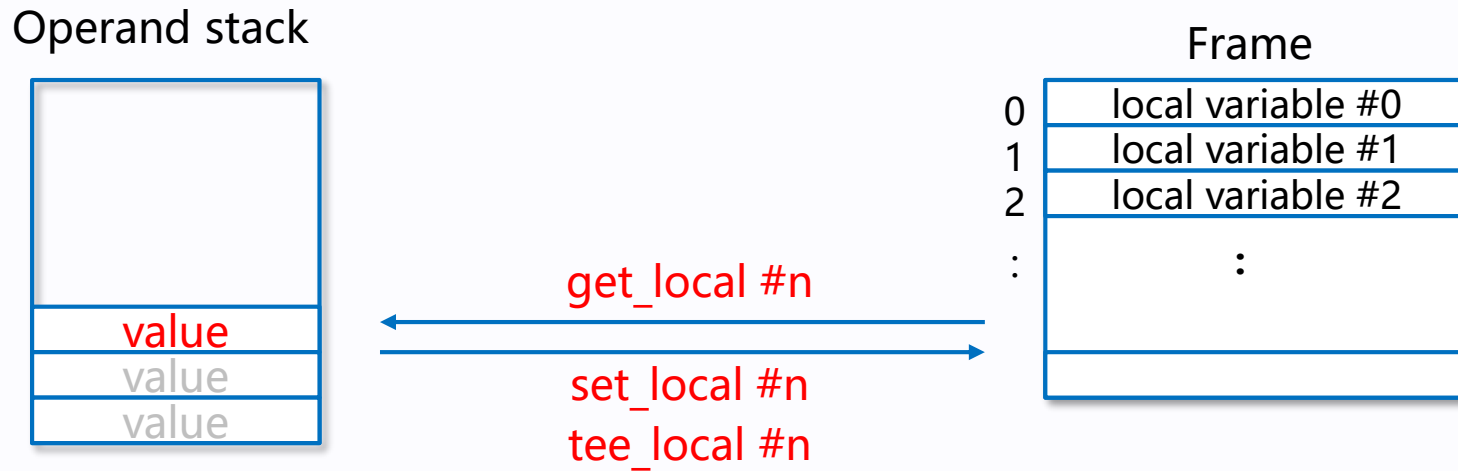
The `select` instruction selects one of its first two operands based on whether its third operand is zero or not.

Global variable instructions



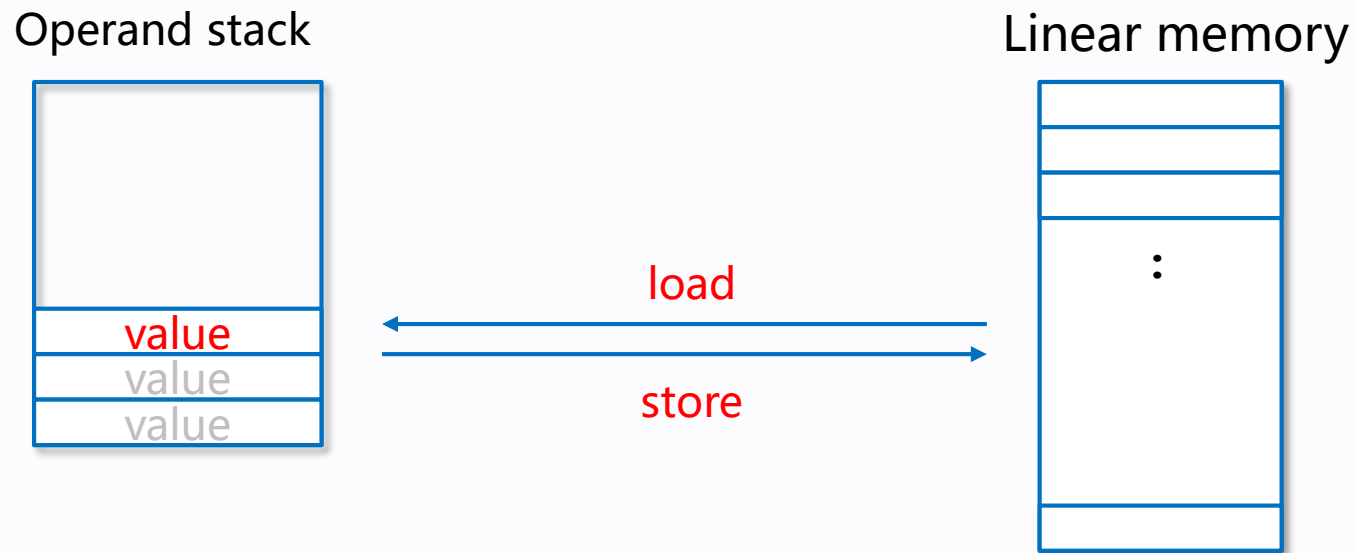
Global variable instructions get or set the values of variables.

Local variable instructions



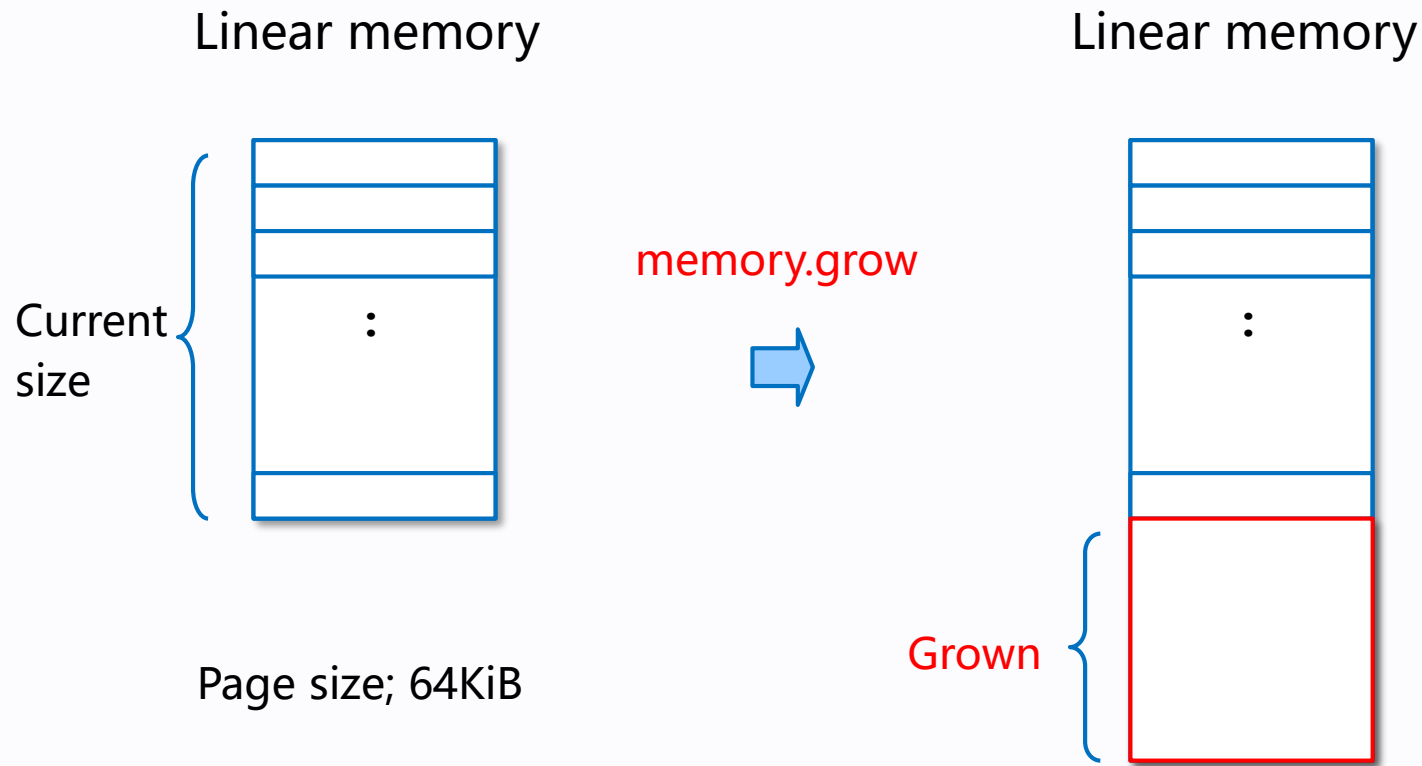
Local variable instructions get or set the values of variables.
(including function arguments)

Memory instructions : load, store



Memory is accessed with load and store instructions for the different value types.

Memory instructions : memory.grow

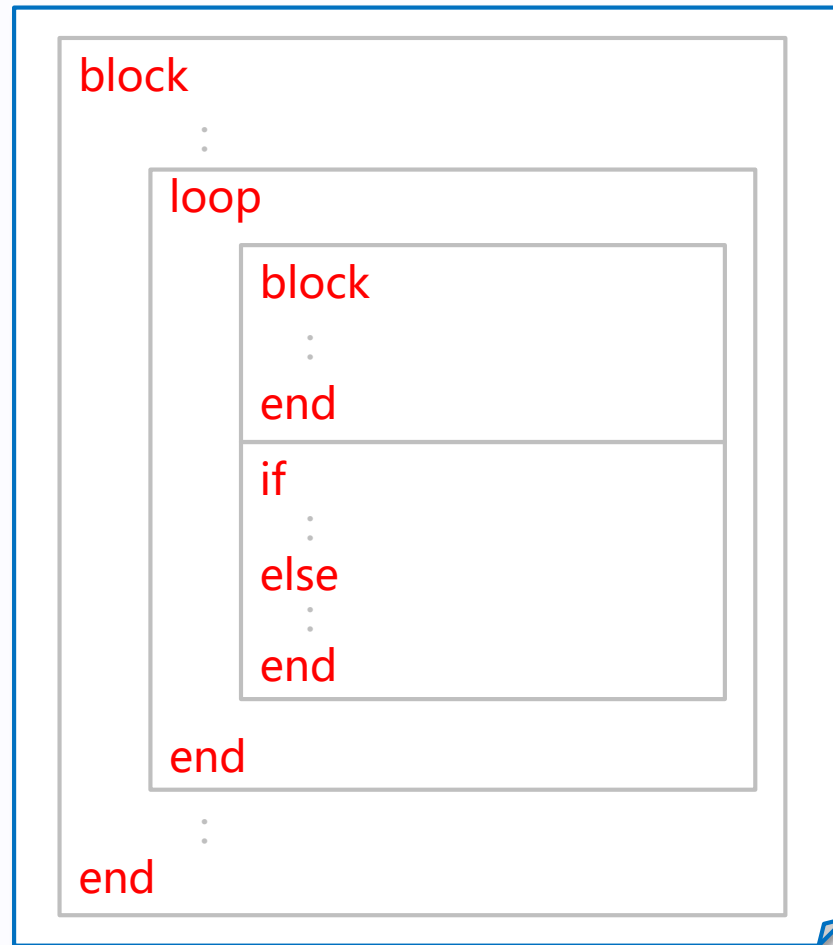


The memory.grow instruction grows memory by a given delta.
The memory.grow instruction operate in units of page size (64KiB).

4. WebAssembly instructions

Control instructions

Control flow is structured

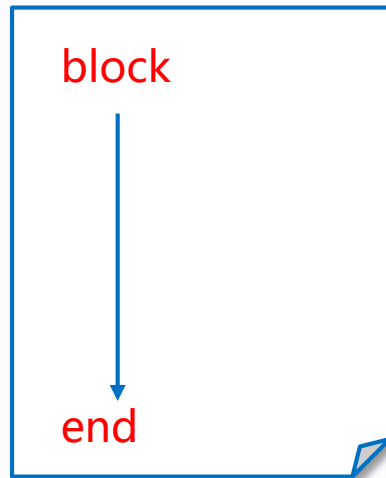


Control flow is expressed with well-nested constructs such as blocks, loops, and conditionals (if-else).

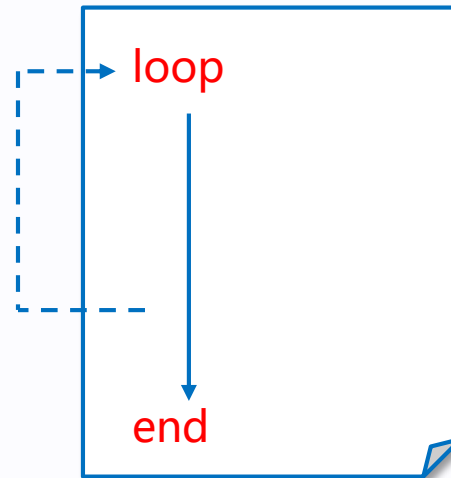
Structured control flow allows simpler and more efficient verification.

Structured control instructions

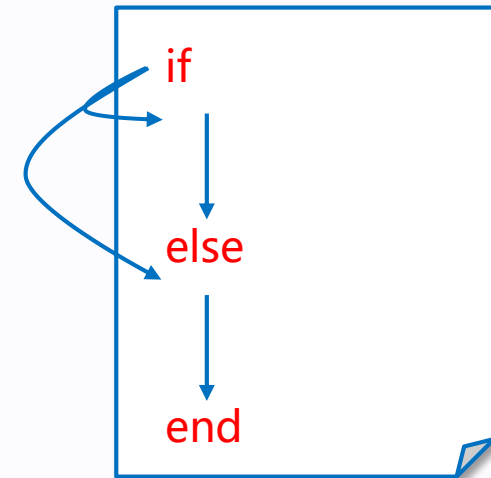
block construct



loop construct



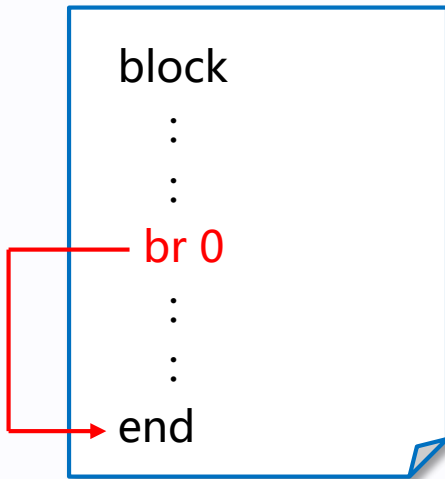
if construct



The block, loop and if instructions are structured control instructions.

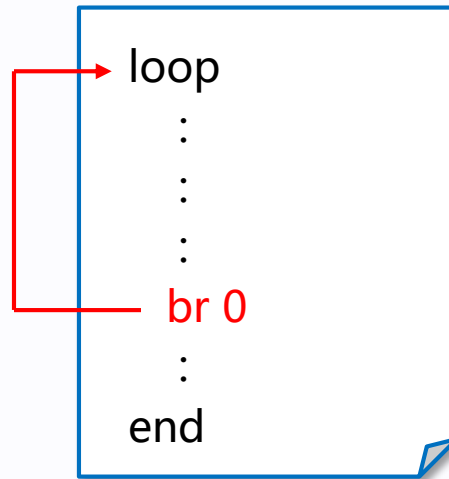
Control constructs and branch instruction

block construct



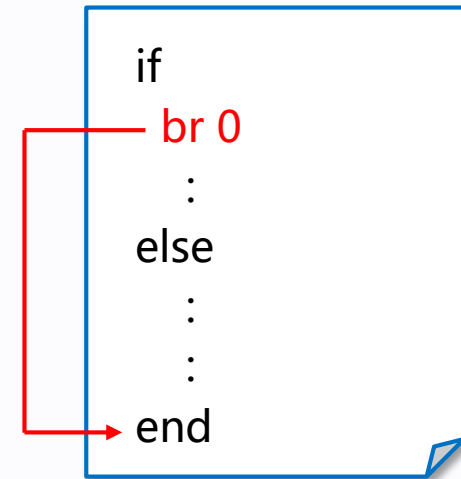
forward jump

loop construct



backward jump

if construct

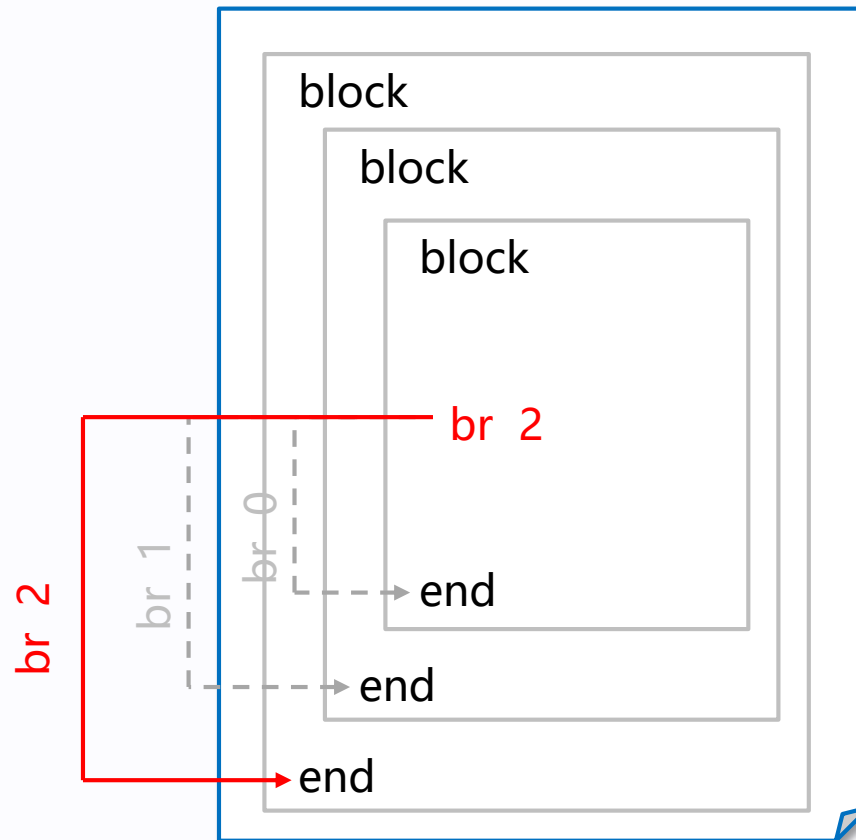


forward jump

Branches can only target control constructs.

Intuitively, a branch targeting a block or if behaves like a break statement, while a branch targeting a loop behaves like a continue statement.

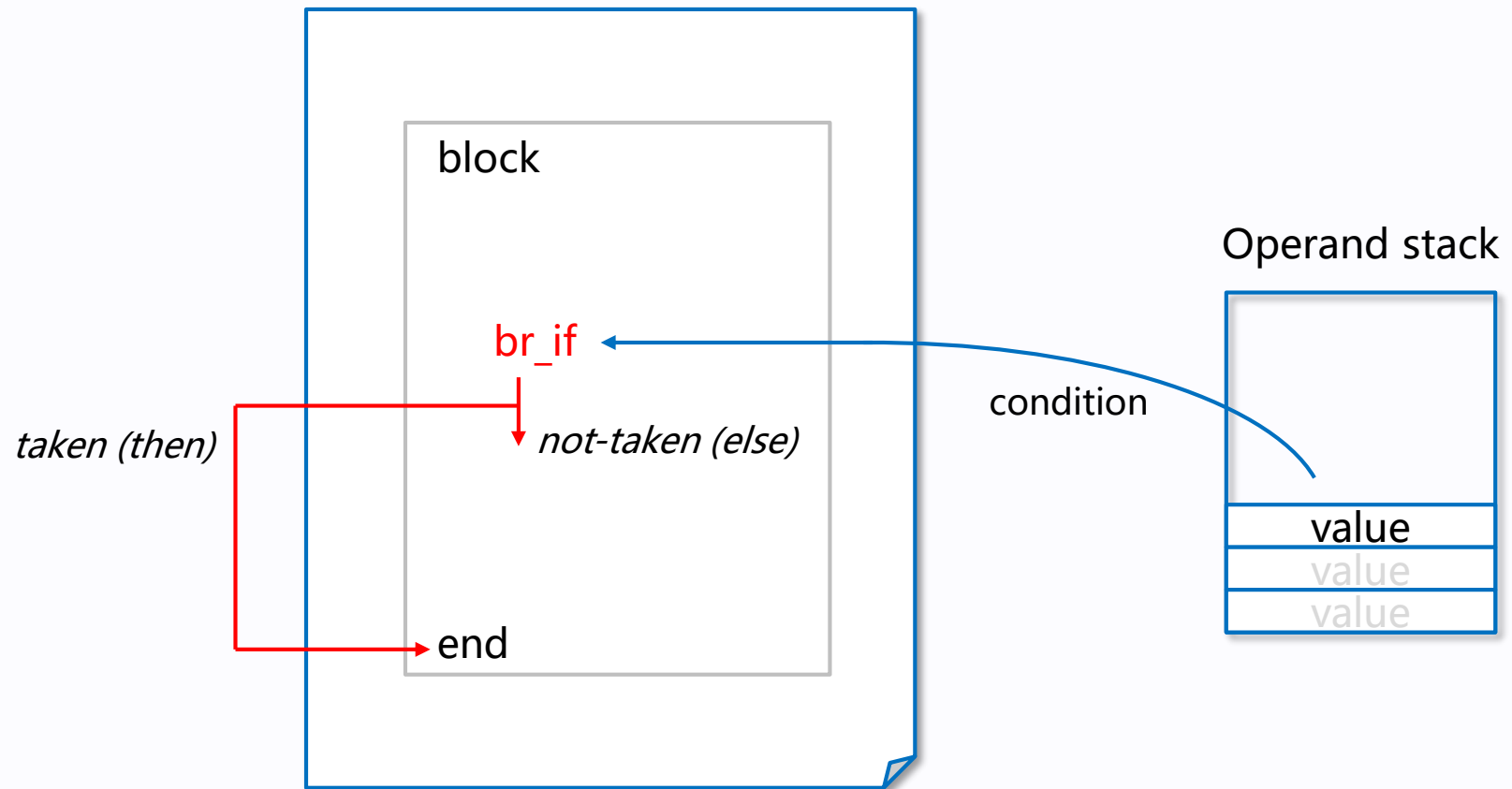
Nested constructs and branch instruction



Branches have "label" immediates.

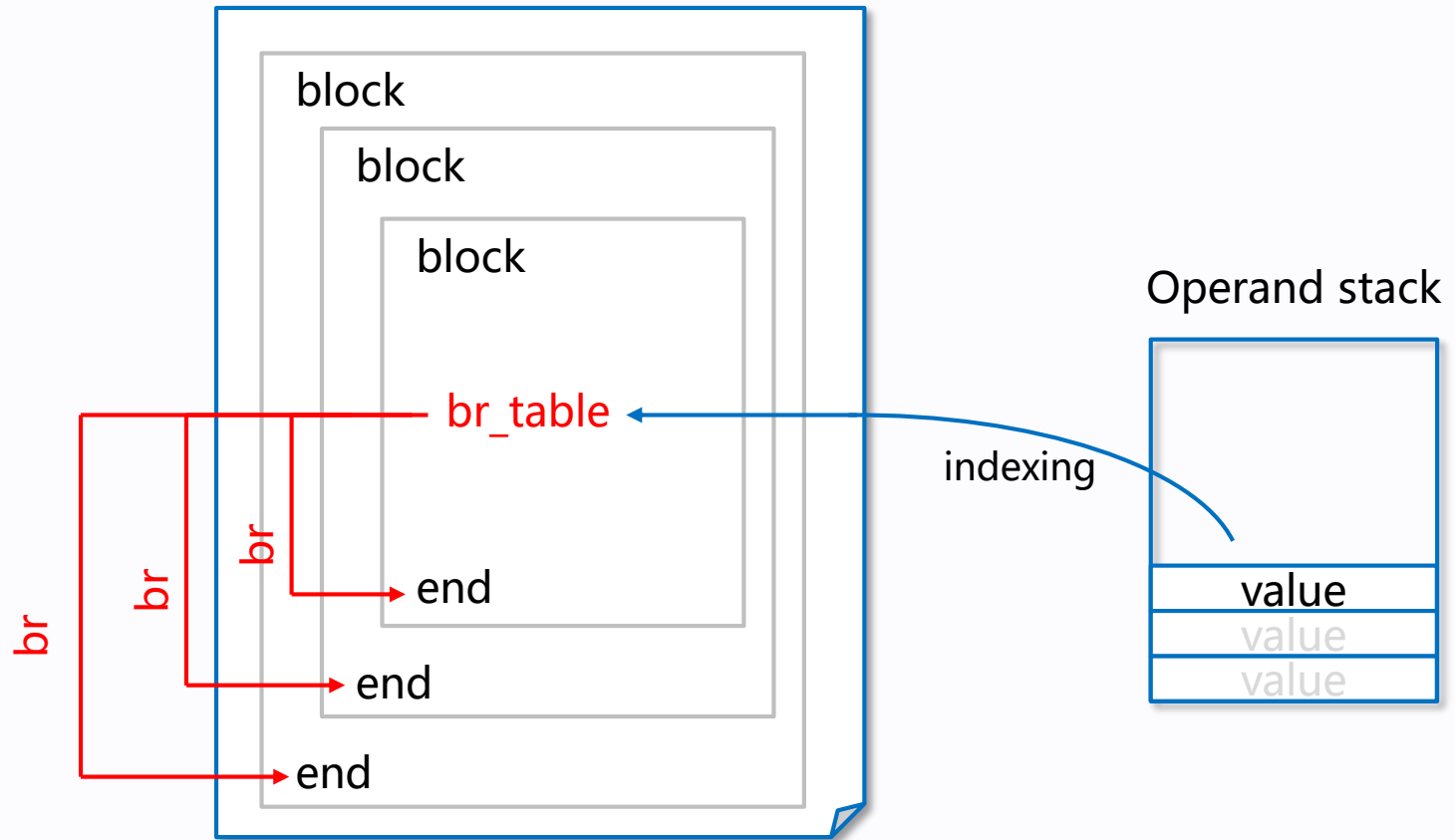
It do not reference program positions in the instruction stream but instead reference outer control constructs by relative nesting depth.

Conditional branch instruction



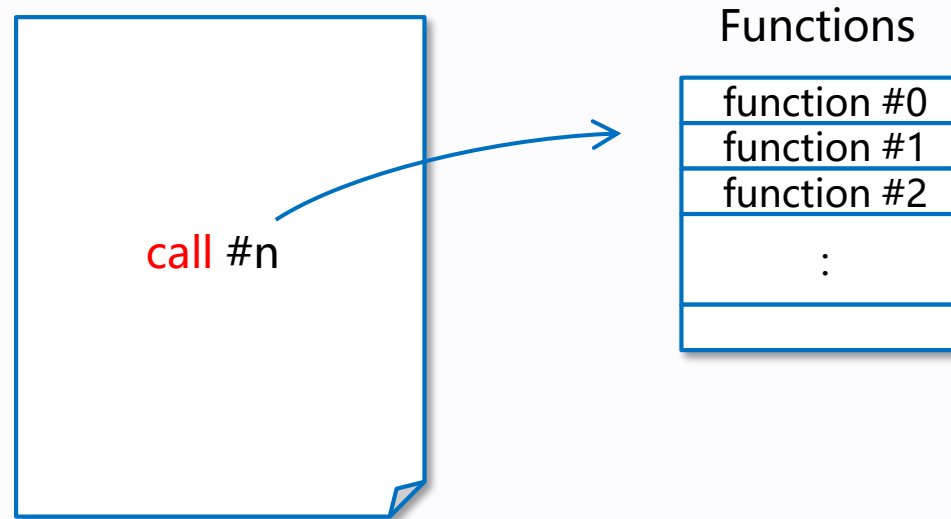
The `br_if` instruction performs a conditional branch.

Table branch instruction



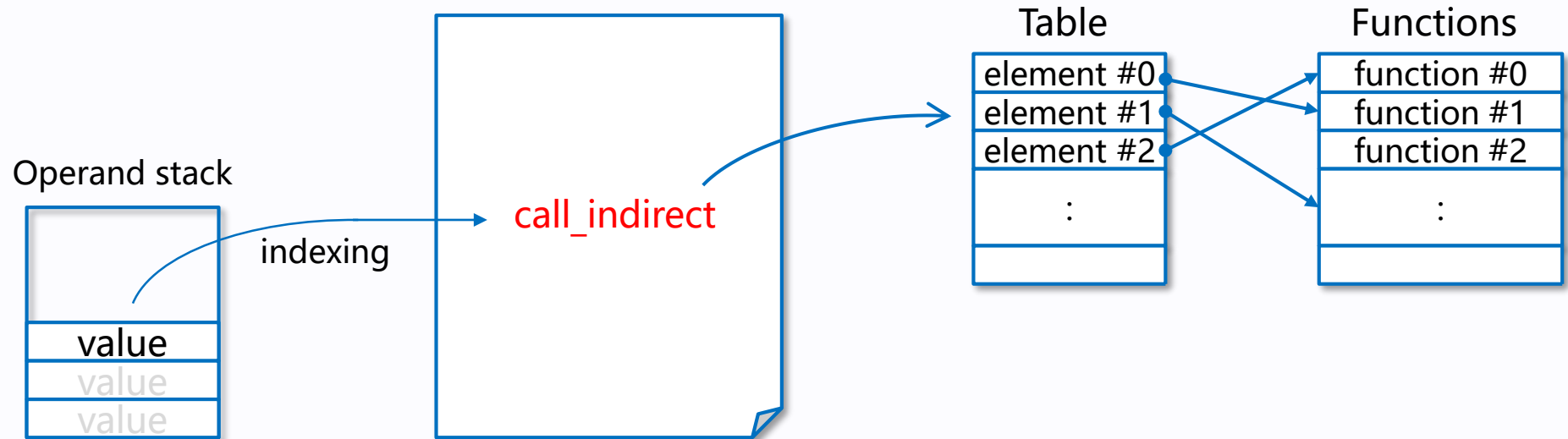
The `br_table` performs an indirect branch through an operand indexing into the label vector.

Call instruction



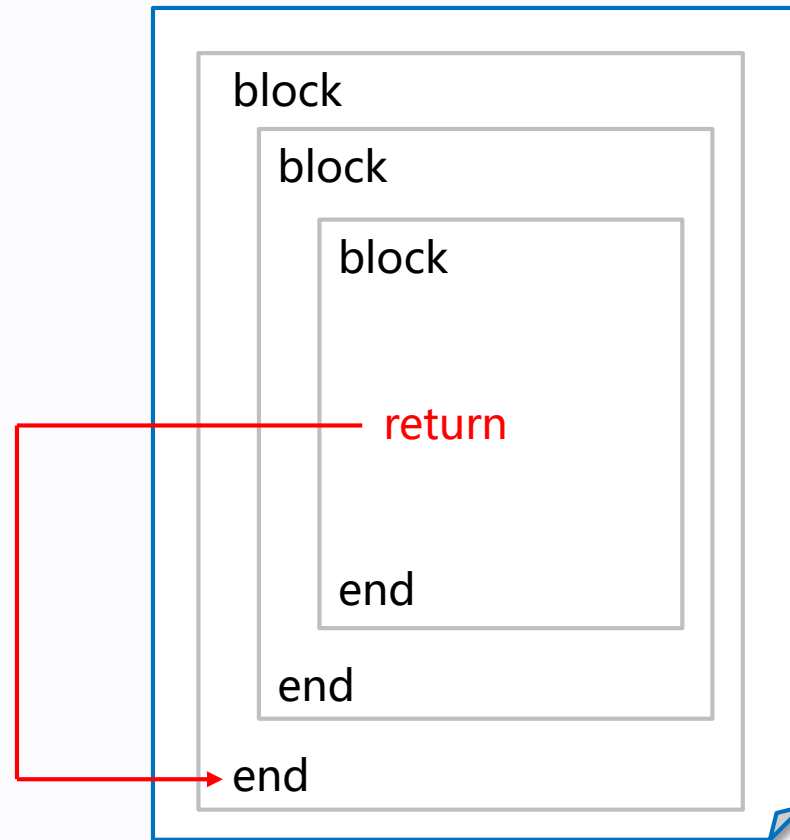
The call instruction invokes another function, consuming the necessary arguments from the stack and returning the result values of the call.

Indirect call instruction



The `call_indirect` instruction calls a function indirectly through an operand indexing into a table.

Return instruction

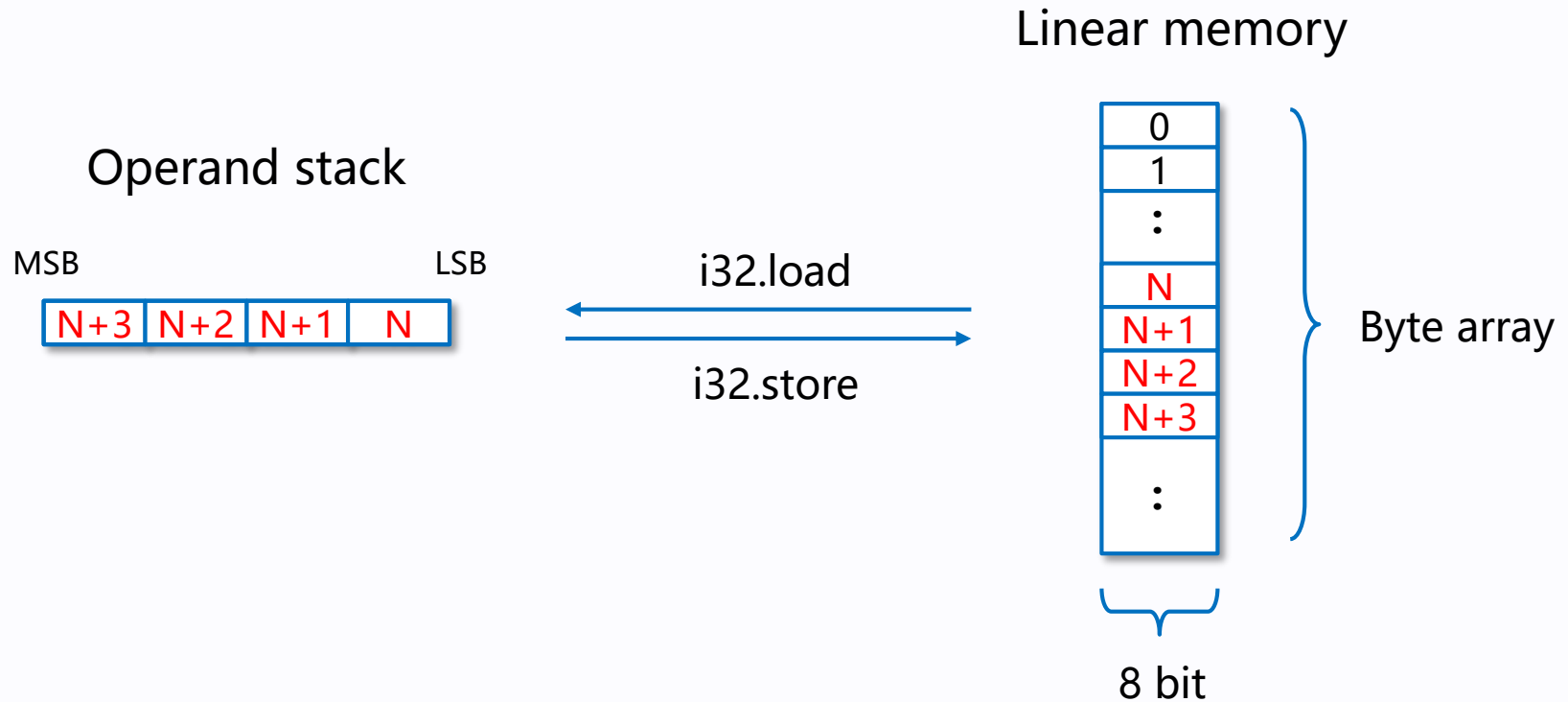


The return instruction is an unconditional branch to the outermost block, which implicitly is the body of the current function.

4. WebAssembly instructions

Byte order

Endian



WebAssembly abstract machine is little endian byte order.
When a number is stored into memory, it is converted into a sequence of bytes in little endian byte order.

Appendix A

Appendix A

Semantics

Validation and execution semantics

The semantics is derived from the following article:
 "Bringing the Web up to Speed with WebAssembly" [2]

Validation semantics: typing rules

Typing Instructions $C \vdash e^* : tf$

$C \vdash t.\text{const } c : \epsilon \rightarrow t$	$C \vdash t.\text{unop} : t \rightarrow t$	$C \vdash t.\text{binop} : t \rightarrow t$	$C \vdash t.\text{testop} : t \rightarrow i32$	$C \vdash t.\text{relop} : t \rightarrow i32$
$t_1 \neq t_2 \quad \text{sz}^? = \epsilon \Leftrightarrow (t_1 = \text{in} \wedge t_2 = \text{in}' \wedge t_1 < t_2) \vee (t_1 = \text{fn} \wedge t_2 = \text{fn}')$				
$C \vdash t_1.\text{convert } t_2.\text{sz}^? : t_2 \rightarrow t_1$			$C \vdash t_1.\text{reinterpret } t_2 : t_2 \rightarrow t_1$	
$C \vdash \text{unreachable} : t_1^? \rightarrow t_2^?$ $C \vdash \text{nop} : \epsilon \rightarrow \epsilon$ $C \vdash \text{drop} : t \rightarrow \epsilon$ $C \vdash \text{select} : t \rightarrow i32 \rightarrow t$				
$tf = t_1^n \rightarrow t_2^m \quad C, \text{label}(t_2^m) \vdash e^* : tf$ $tf = t_1^n \rightarrow t_2^m \quad C, \text{label}(t_1^n) \vdash e^* : tf$				
$C \vdash \text{block } tf \ e^* \ \text{end} : tf$ $C \vdash \text{loop } tf \ e^* \ \text{end} : tf$				
$tf = t_1^n \rightarrow t_2^m \quad C, \text{label}(t_2^m) \vdash e_1^* : tf$ $C, \text{label}(t_2^m) \vdash e_2^* : tf$				
$C \vdash \text{if } tf \ e_1^* \ \text{else } e_2^* \ \text{end} : t_1^? \rightarrow t_2^?$				
$C_{\text{label}}(i) = t^*$ $C_{\text{label}}(i) = t^*$ $(C_{\text{label}}(i) = t^*)^+$				
$C \vdash \text{br } i : t_1^? \rightarrow t_2^?$ $C \vdash \text{br } i : t^* \rightarrow t^*$ $C \vdash \text{br } i : t^* \rightarrow t^*$ $C \vdash \text{br } i : t^* \rightarrow t^*$				
$C_{\text{return}} = t^*$ $C_{\text{unc}}(i) = tf$ $tf = t_1^? \rightarrow t_2^?$ $C_{\text{table}} = n$				
$C \vdash \text{return} : t_1^? \rightarrow t_2^?$ $C \vdash \text{call } i : tf$ $C \vdash \text{call } i : tf$ $C \vdash \text{call } i : tf$				
$C_{\text{local}}(i) = t$ $C_{\text{local}}(i) = t$ $C_{\text{local}}(i) = t$ $C_{\text{global}}(i) = \text{mut}^? t$ $C_{\text{global}}(i) = \text{mut } t$				
$C \vdash \text{get } i : \epsilon \rightarrow t$ $C \vdash \text{set } i : t \rightarrow \epsilon$ $C \vdash \text{tee } i : t \rightarrow t$ $C \vdash \text{get } i : \epsilon \rightarrow t$ $C \vdash \text{set } i : t \rightarrow \epsilon$				
$C_{\text{memory}} = n \quad 2^* \leq (tp < t) \quad (tp.\text{sz}^?) = \epsilon \vee t = \text{im}$ $C_{\text{memory}} = n \quad 2^* \leq (tp < t) \quad tp^? = \epsilon \vee t = \text{im}$				
$C \vdash t.\text{load } (tp.\text{sz}^?) \ a \ o : i32 \rightarrow t$ $C \vdash t.\text{store } tp^? \ a \ o : i32 \rightarrow \epsilon$				
$C_{\text{memory}} = n$ $C_{\text{memory}} = n$				
$C \vdash \text{current } \text{memory} : \epsilon \rightarrow i32$ $C \vdash \text{grow } \text{memory} : i32 \rightarrow i32$				
$C \vdash \epsilon : \epsilon \rightarrow \epsilon$ $C \vdash e_1^* : t_1^? \rightarrow t_2^?$ $C \vdash e_2^* : t_2^? \rightarrow t_3^?$ $C \vdash e^* : t_1^? \rightarrow t_2^?$				
$tf = t_1^? \rightarrow t_2^? \quad C, \text{local } t_1^? \ t_2^?, \text{label}(t_2^?), \text{return}(t_2^?) \vdash e^* : \epsilon \rightarrow t_2^?$	$tg = \text{mut}^? t \quad C \vdash e^* : \epsilon \rightarrow t \quad \text{ex}^* = \epsilon \vee tg = t$			
$C \vdash \text{ex}^* \ \text{func } tf \ \text{local } t^* \ e^* : \text{ex}^* \ tf$	$C \vdash \text{ex}^* \ \text{global } tg \ e^* : \text{ex}^* \ tg$			
$(C_{\text{unc}}(i) = tf)^n$				
$C \vdash \text{ex}^* \ \text{table } n \ i^n : \text{ex}^* \ n$ $C \vdash \text{ex}^* \ \text{memory } n : \text{ex}^* \ n$				
$tg = t$				
$C \vdash \text{ex}^* \ \text{func } tf \ \text{im} : \text{ex}^* \ tf$	$C \vdash \text{ex}^* \ \text{global } tg \ \text{im} : \text{ex}^* \ tg$	$C \vdash \text{ex}^* \ \text{table } n \ \text{im} : \text{ex}^* \ n$	$C \vdash \text{ex}^* \ \text{memory } n \ \text{im} : \text{ex}^* \ n$	
$(C \vdash f : \text{ex}_f^? \ tf)^*$ $(C \vdash \text{glob}_g : \text{ex}_g^? \ tg)_i^*$ $(C \vdash \text{tab}_n : \text{ex}_n^? \ n)^?$ $(C \vdash \text{mem} : \text{ex}_m^? \ n)^?$				
$(C_s = \{\text{global } tg^{i-1}\})_i^*$ $C = \{\text{func } tf^*, \text{global } tg^*, \text{table } n^*, \text{memory } n^*\}$ $\text{ex}_i^* \ \text{ex}_g^* \ \text{ex}_n^* \ \text{ex}_m^* \ \text{distinct}$				
$\vdash \text{module } f^* \ \text{glob}^* \ \text{tab}^? \ \text{mem}^?$				

Execution semantics: reduction rules

Reduction $s; v^*; e^* \hookrightarrow_t s'; v'^*; e'^*$

$s; v^*; L^k[e^*] \hookrightarrow_s s'; v'^*; L^k[e'^*]$	$s; v^*; \text{local}_n(i; v^*) \ e^* \ \text{end} \hookrightarrow_{s'} s'; v'^*; \text{local}_n(i; v'^*) \ e'^* \ \text{end}$
$L^0[\text{trap}] \hookrightarrow \text{trap}$ if $L^0 \neq []$	
$(t.\text{const } c) \ t.\text{unop} \hookrightarrow t.\text{const } \text{unop}_1(c)$	
$(t.\text{const } c_1) \ (t.\text{const } c_2) \ t.\text{binop} \hookrightarrow t.\text{const } c$ if $c = \text{binop}_1(c_1, c_2)$	
$(t.\text{const } c_1) \ (t.\text{const } c_2) \ t.\text{binop} \hookrightarrow \text{trap}$ otherwise	
$(t.\text{const } c_1) \ (t.\text{const } c_2) \ t.\text{testop} \hookrightarrow i32.\text{const } \text{testop}_1(c)$	
$(t.\text{const } c_1) \ (t.\text{const } c_2) \ t.\text{relop} \hookrightarrow i32.\text{const } \text{relop}_1(c_1, c_2)$	
$(t_1.\text{const } c) \ t_2.\text{convert } t_1.\text{sz}^? \hookrightarrow t_2.\text{const } c'$ if $c' = \text{cv}_{t_1, t_2}^{\text{sz}^?}(c)$	
$(t_1.\text{const } c) \ t_2.\text{convert } t_1.\text{sz}^? \hookrightarrow \text{trap}$ otherwise	
$(t_1.\text{const } c) \ t_2.\text{reinterpret } t_1 \hookrightarrow t_2.\text{const } \text{const}_{e_2}(\text{bits}_{t_1}(c))$	
unreachable $\hookrightarrow \text{trap}$	
nop $\hookrightarrow \epsilon$	
v drop $\hookrightarrow \epsilon$	
$v_1 \ v_2 \ (i32.\text{const } 0) \ \text{select} \hookrightarrow v_2$	
$v_1 \ v_2 \ (i32.\text{const } k + 1) \ \text{select} \hookrightarrow v_1$	
$v^n \ \text{block } (t_1^n \rightarrow t_2^m) \ e^* \ \text{end} \hookrightarrow \text{label}_n[\epsilon] \ v^n \ e^* \ \text{end}$	
$v^n \ \text{loop } (t_1^n \rightarrow t_2^m) \ e^* \ \text{end} \hookrightarrow \text{label}_n[\text{loop } (t_1^n \rightarrow t_2^m) \ e^* \ \text{end}] \ v^n \ e^* \ \text{end}$	
$(i32.\text{const } 0) \ \text{if } tf \ e_1^* \ \text{else } e_2^* \ \text{end} \hookrightarrow \text{block } tf \ e_2^* \ \text{end}$	
$(i32.\text{const } k + 1) \ \text{if } tf \ e_1^* \ \text{else } e_2^* \ \text{end} \hookrightarrow \text{block } tf \ e_1^* \ \text{end}$	
$\text{label}_n[e^*] \ v^* \ \text{end} \hookrightarrow v^*$	
$\text{label}_n[e^*] \ \text{trap} \ \text{end} \hookrightarrow \text{trap}$	
$\text{label}_n[e^*] \ L^? [v^n \ (\text{br } j)] \ \text{end} \hookrightarrow v^n \ e^*$	
$(i32.\text{const } 0) \ (\text{br } j) \hookrightarrow \epsilon$	
$(i32.\text{const } k + 1) \ (\text{br } j) \hookrightarrow \text{br } j$	
$(i32.\text{const } k) \ (\text{br } \text{table } j_1^? \ j_2^?) \hookrightarrow \text{br } j$	
$(i32.\text{const } k + n) \ (\text{br } \text{table } j_1^? \ j) \hookrightarrow \text{br } j$	
$s; \text{call } j \hookrightarrow \text{call } s_{\text{unc}}(i, j)$	
$s; (i32.\text{const } j) \ \text{call } \text{indirect } tf \hookrightarrow \text{call } s_{\text{tab}}(i, j)$ if $s_{\text{tab}}(i, j)_{\text{code}} = (\text{func } tf \ \text{local } t^* \ e^*)$	
$s; (i32.\text{const } j) \ \text{call } \text{indirect } tf \hookrightarrow \text{trap}$ otherwise	
$v^n \ (\text{call } cl) \hookrightarrow \text{local}_m[cl_{\text{mem}}; v^n \ (t.\text{const } 0)^k] \ \text{block } (\epsilon \rightarrow t_2^m) \ e^* \ \text{end} \ \text{end} \ \dots$	
$\text{local}_n[i; v_i^*] \ v^n \ \text{end} \hookrightarrow v^n$... if $cl_{\text{code}} = (\text{func } (t_1^n \rightarrow t_2^m) \ \text{local } t^k \ e^*)$	
$\text{local}_n[i; v_i^*] \ L^? [v^n \ \text{return}] \ \text{end} \hookrightarrow v^n$	
$v_j^? \ v \ v_2^?; \text{get } \text{local } j \hookrightarrow v_j^? \ v \ v_2^?; \epsilon$	
$v_j^? \ v \ v_2^?; v' \ (\text{set } \text{local } j) \hookrightarrow v \ v \ (\text{set } \text{local } j)$	
$v \ (\text{tee } \text{local } j) \hookrightarrow v \ v \ (\text{tee } \text{local } j)$	
$s; \text{get } \text{global } j \hookrightarrow s_{\text{glob}}(i, j)$	
$s; v \ (\text{set } \text{global } j) \hookrightarrow s'; \epsilon$ if $s' = s$ with $\text{glob}(i, j) = v$	
$s; (i32.\text{const } k) \ (t.\text{load } a \ o) \hookrightarrow t.\text{const } \text{const}_{t_1}(b^*)$ if $s_{\text{mem}}(i, k + o, t) = b^*$	
$s; (i32.\text{const } k) \ (t.\text{load } tp.\text{sz} \ a \ o) \hookrightarrow t.\text{const } \text{const}_{t_1}^{\text{sz}^?}(b^*)$ if $s_{\text{mem}}(i, k + o, tp) = b^*$	
$s; (i32.\text{const } k) \ (t.\text{load } tp.\text{sz}^? \ a \ o) \hookrightarrow \text{trap}$ otherwise	
$s; (i32.\text{const } k) \ (t.\text{const } c) \ (t.\text{store } a \ o) \hookrightarrow s'; \epsilon$ if $s' = s$ with $\text{mem}(i, k + o, t) = \text{bits}_{t_1}^{ t }(c)$	
$s; (i32.\text{const } k) \ (t.\text{const } c) \ (t.\text{store } tp \ a \ o) \hookrightarrow s'; \epsilon$ if $s' = s$ with $\text{mem}(i, k + o, tp) = \text{bits}_{tp}^{ tp }(c)$	
$s; (i32.\text{const } k) \ (t.\text{const } c) \ (t.\text{store } tp^? \ a \ o) \hookrightarrow \text{trap}$ otherwise	
$s; \text{current } \text{memory} \hookrightarrow i32.\text{const } \lfloor s_{\text{mem}}(i, *) \rfloor / 64 \text{ Ki}$	
$s; (i32.\text{const } k) \ \text{grow } \text{memory} \hookrightarrow s'; i32.\text{const } \lfloor s_{\text{mem}}(i, *) \rfloor / 64 \text{ Ki}$ if $s' = s$ with $\text{mem}(i, *) = s_{\text{mem}}(i, *) \ (0)^{k-64 \text{ Ki}}$	
$s; (i32.\text{const } k) \ \text{grow } \text{memory} \hookrightarrow i32.\text{const } (-1)$	

Figure 2. Small-step reduction rules

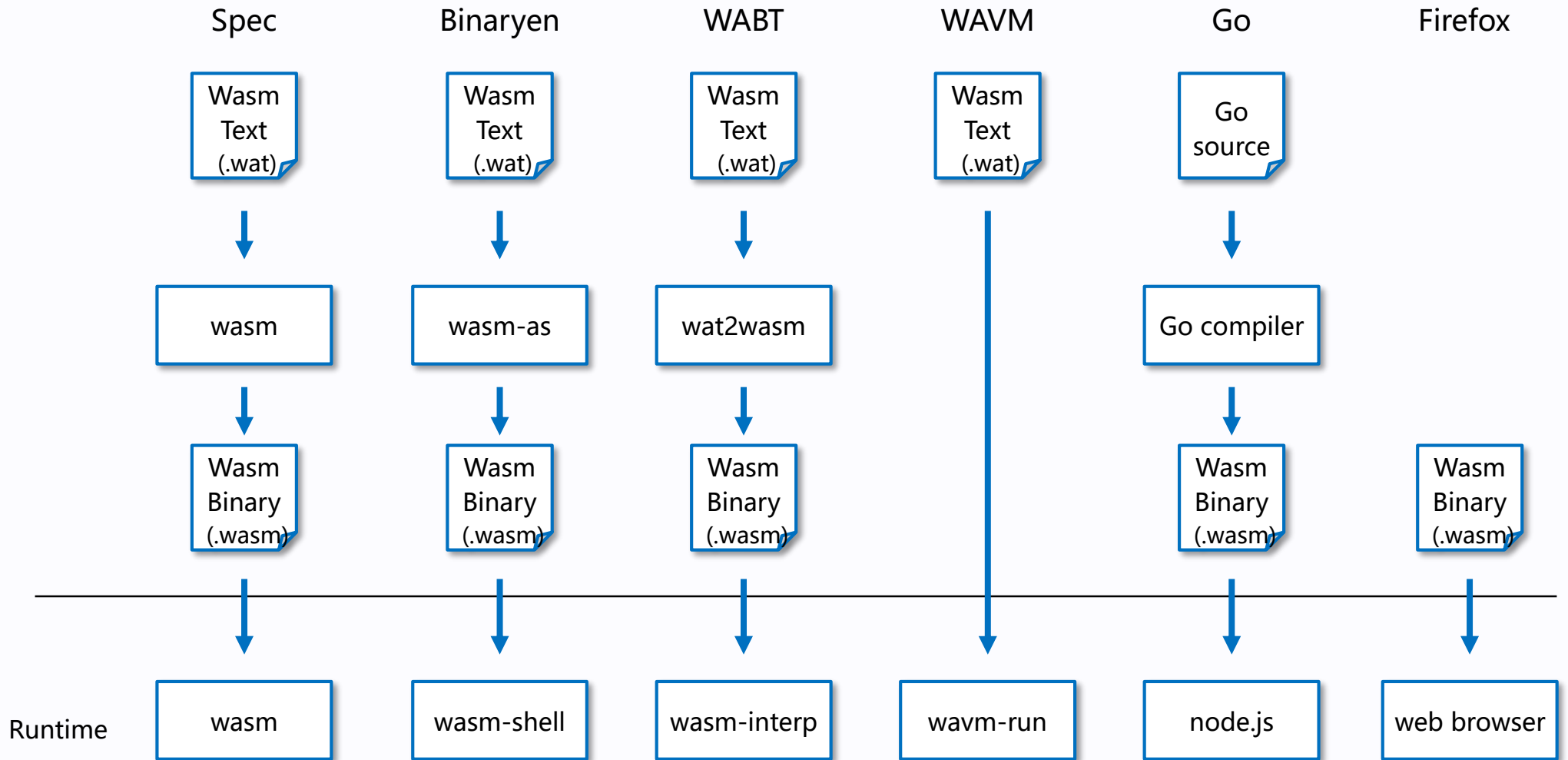
References : [2], [1] Ch.3, Ch.4, Ch.7

Appendix B

Appendix B

Implementations

Implementations



Reference interpreter : spec

<https://github.com/WebAssembly/spec>
[interpreter/exec/eval.ml]

```
let rec step (c : config) : config =
  let {frame; code = vs, es; _} = c in
  let e = List.hd es in
  let vs', es' =
    match e.it, vs with
    | Plain e', vs ->
      (match e', vs with
      | Unreachable, vs ->
        vs, [Trapping "unreachable executed" @@ e.at]

      | Nop, vs ->
        vs, []

      | Block (ts, es'), vs ->
        vs, [Label (List.length ts, [], ([], List.map plain es')) @@ e.at]

      | Loop (ts, es'), vs ->
        vs, [Label (0, [e' @@ e.at], ([], List.map plain es')) @@ e.at]

      | If (ts, es1, es2), I32 01 :: vs' ->
        vs', [Plain (Block (ts, es2)) @@ e.at]
```

Interpreter : WABT

<https://github.com/WebAssembly/wabt>

[src/interp/interp.cc]

```
Result Thread::Run(int num_instructions) {
    Result result = Result::Ok;

    const uint8_t* istream = GetIstream();
    const uint8_t* pc = &istream[pc_];
    for (int i = 0; i < num_instructions; ++i) {
        Opcode opcode = ReadOpcode(&pc);
        assert(!opcode.IsInvalid());
        switch (opcode) {
            case Opcode::Select: {
                uint32_t cond = Pop<uint32_t>();
                Value false_ = Pop();
                Value true_ = Pop();
                CHECK_TRAP(Push(cond ? true_ : false_));
                break;
            }

            case Opcode::Br:
                GOTO(ReadU32(&pc));
                break;

            case Opcode::BrIf: {
```

Stand-alone VM : WAVM

<https://github.com/WAVM/WAVM>

[Lib/LLVMJIT/EmitFunction.cpp]

```
// Decode the WebAssembly opcodes and emit LLVM IR for them.
    OperatorDecoderStream decoder(functionDef.code);
    UnreachableOpVisitor unreachableOpVisitor(*this);
    OperatorPrinter operatorPrinter(irModule, functionDef);
    Uptr opIndex = 0;
    while(decoder && controlStack.size())
    {
        irBuilder.SetCurrentDebugLocation(
            llvm::DILocation::get(llvmContext, (unsigned int)opIndex++,
0, diFunction));
        if(ENABLE_LOGGING)
{ logOperator(decoder.decodeOpWithoutConsume(operatorPrinter)); }

        if(controlStack.back().isReachable) { decoder.decodeOp(*this); }
        else
        {
            decoder.decodeOp(unreachableOpVisitor);
        }
        wvmAssert(irBuilder.GetInsertBlock() == returnBlock);

        if(EMIT_ENTER_EXIT_HOOKS)
```

Web browser : Firefox

<https://github.com/mozilla/gecko-dev>
[js/src/wasm/WasmBaselineCompile.cpp]

```
switch (op.b0) {
    case uint16_t(Op::End):
        if (!emitEnd()) {
            return false;
        }

        if (iter_.controlStackEmpty()) {
            if (!deadCode_) {
                doReturn(funcType().ret(), PopStack(false));
            }
            return iter_.readFunctionEnd(iter_.end());
        }
        NEXT();

    // Control opcodes
    case uint16_t(Op::Nop):
        CHECK_NEXT(iter_.readNop());
    case uint16_t(Op::Drop):
        CHECK_NEXT(emitDrop());
    case uint16_t(Op::Block):
        CHECK_NEXT(emitBlock());
    case uint16_t(Op::Loop):
```

Appendix B

CLI development utilities

Assemble

Assemble Wasm text format (.wat) to Wasm binary format (.wasm) :

Binaryen :

```
$ wasm-as sample.wat
```

WABT :

```
$ wat2wasm sample.wat
```

```
$ wat2wasm -v sample.wat
```

Spec :

```
$ wasm -d sample.wat
```


Disassemble

Disassemble Wasm binary format (.wasm) to Wasm text format (.wat)

Binaryen :

```
$ wasm-dis sample.wasm
```

WABT :

```
$ wasm2wat sample.wasm
```

```
$ wasm-objdump -d sample.wasm
```

Spec :

```
$ wasm -d sample.wasm
```

Desugar

Desugar Wasm text format (.wat) to Wasm text format (.wat)

WABT :

```
$ wat-desugar sample.wat
```

Dump information

Dump Wasm binary format (.wasm) information :

WABT :

```
$ wasm-objdump -s sample.wasm
```

```
$ wasm-objdump -x sample.wasm
```

Spec :

```
$ wasm -s sample.wasm
```

Run

Run Wasm binary format (.wasm) and Wasm text format (.wat) :

WABT : Run Wasm binary format with trace

```
$ wasm-interp --run-all-exports --trace sample.wasm
```

WAVM : Run Wasm text format

```
$ wavm-run sample.wat
```

Spec : Run Wasm binary format

```
$ wasm sample.wasm -e '(invoke "XXX")'
```

REPL

REPL (Read-Eval-Print-Loop) :

Spec :

```
$ wasm -
```

```
$ wasm sample.wasm -
```

Appendix B

Test suite

Test suite and Wasm text format examples

<https://github.com/WebAssembly/spec>
[test/core]

```
README.md          fac.wast           names.wast
address.wast       float_exprs.wast  nop.wast
align.wast         float_literals.wast return.wast
binary.wast        float_memory.wast run.py*
block.wast         float_misc.wast   select.wast
br.wast            forward.wast      set_local.wast
br_if.wast         func.wast         skip-stack-guard-page.wast
br_table.wast     func_ptrs.wast   stack.wast
break-drop.wast   get_local.wast   start.wast
call.wast          globals.wast     store_retval.wast
call_indirect.wast i32.wast         switch.wast
comments.wast     i64.wast         tee_local.wast
const.wast        if.wast          token.wast
conversions.wast  imports.wast     traps.wast
custom.wast       inline-module.wast type.wast
data.wast          int_exprs.wast   typecheck.wast
elem.wast          int_literals.wast unreachable.wast
endianness.wast   labels.wast      unreached-invalid.wast
exports.wast      left-to-right.wast unwind.wast
f32.wast           linking.wast     utf8-custom-section-id.wast
f32_bitwise.wast  loop.wast        utf8-import-field.wast
f32_cmp.wast      memory.wast      utf8-import-module.wast
f64.wast           memory_grow.wast utf8-invalid-encoding.wast
f64_bitwise.wast  memory_redundancy.wast
f64_cmp.wast      memory_trap.wast
```

Note: `.wast` extension means command-script and Wasm text format.

References: [C1]

Appendix B

Desugar examples

Desugar example

Text format
syntactic sugar

```
(func (result i32)
  (i32.add
    (i32.const 1)
    (i32.const 2)))
```



Text format
core syntax

```
(func (result i32)
  i32.const 1
  i32.const 2
  i32.add)
```

Desugar example

Text format
syntactic sugar

```
(func (result i32)
  (i32.add
    (i32.const 1)
    (i32.mul
      (i32.const 2)
      (i32.const 3))))
```



Text format
core syntax

```
(func (result i32)
  i32.const 1
  i32.const 2
  i32.const 3
  i32.mul
  i32.add)
```

Desugar example

Text format
syntactic sugar

```
(func (result i32)
  (block (result i32)
    (i32.add
      (i32.const 1)
      (i32.const 2))))
```



Text format
core syntax

```
(func (result i32)
  block (result i32)
    i32.const 1
    i32.const 2
    i32.add
  end)
```

Desugar example

Text format
syntactic sugar

```
(func
  (block $label_a
    (block $label_b
      br $label_a)))
```



Text format
core syntax

```
(func
  block
  block
  br 1
  end
end)
```

Desugar example

Text format
syntactic sugar

```
(func (result i32)
  (if (result i32) (get_global 0)
    (then (i32.const 1))
    (else (i32.const 2))))
```



Text format
core syntax

```
(func (result i32)
  get_global 0
  if (result i32)
    i32.const 1
  else
    i32.const 2
  end)
```

Appendix C

Future

Future directions

- * zero-cost exception, threads, SIMD
- * tail call, stack switching, coroutines
- * garbage collectors

References

References

- [1] WebAssembly Specification Release 1.0 (Draft, last updated Oct 31, 2018)
<https://webassembly.github.io/spec/core/>
- [2] Bringing the Web up to Speed with WebAssembly
<https://github.com/WebAssembly/spec/blob/master/papers/pldi2017.pdf>
- [3] WebAssembly High-Level Goals
<https://webassembly.org/docs/high-level-goals/>
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<https://webassembly.org/docs/rationale/>
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https://developer.mozilla.org/en-US/docs/WebAssembly/Understanding_the_text_format
- [8] Wikipedia: LEB128
<https://en.wikipedia.org/wiki/LEB128>

References

- [C1] spec: WebAssembly specification, reference interpreter, and test suite.
<https://github.com/WebAssembly/spec>
- [C2] Binaryen: Compiler infrastructure and toolchain library for WebAssembly, in C++
<https://github.com/WebAssembly/binaryen>
- [C3] WABT: The WebAssembly Binary Toolkit
<https://github.com/WebAssembly/wabt>
- [C4] WAVM: WebAssembly Virtual Machine
<https://github.com/WAVM/WAVM>
- [C5] mozilla/gecko-dev (Firefox)
<https://github.com/mozilla/gecko-dev>

Here is the slide: <https://github.com/takenobu-hs/WebAssembly-illustrated>